The kernel report

(ELC 2012 edition)

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The Plan

Look at a year's worth of kernel work ...with an eye toward the future



Starting off 2011

2.6.37 released - January 4, 2011 11,446 changes, 1,276 developers

VFS scalability work (inode_lock removal)
Block I/O bandwidth controller
PPTP support
Basic pNFS support
Wakeup sources



What have we done since then?

Since 2.6.37:

Five kernel releases have been made 59,000 changes have been merged 3069 developers have contributed to the kernel 416 companies have supported kernel development



February



As you can see in these posts, Ralink is sending patches for the upstream rt2x00 driver for their new chipsets, and not just dumping a huge, stand-alone tarball driver on the community, as they have done in the past. This shows a huge willingness to learn how to deal with the kernel community, and they should be strongly encouraged and praised for this major change in attitude.

- Greg Kroah-Hartman, February 9

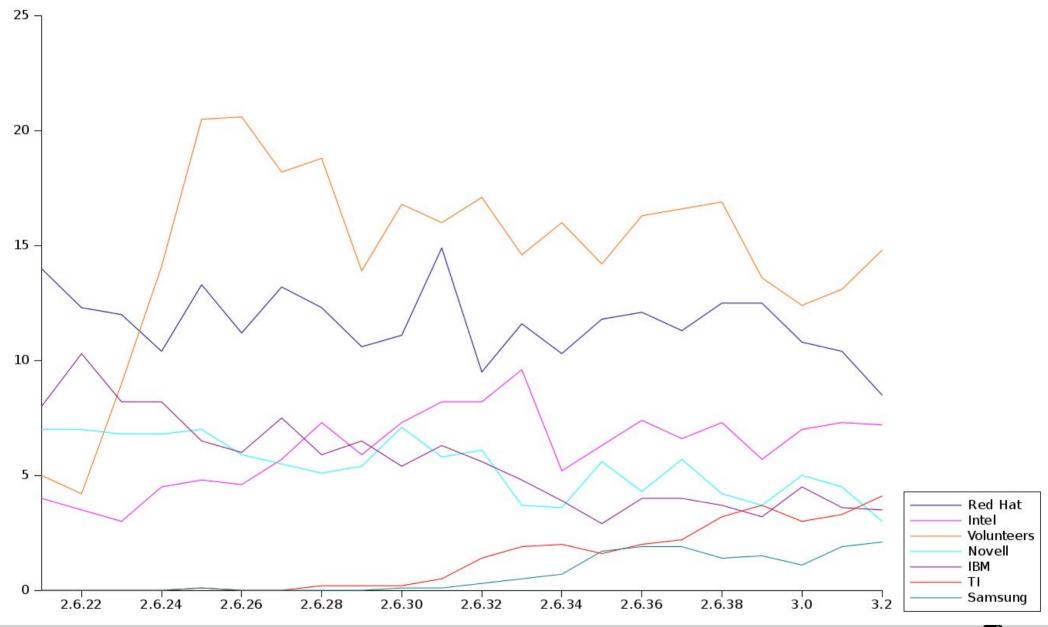


Employer contributions 2.6.38-3.2

Volunteers	13.9%	Wolfson Micro	1.7%
Red Hat	10.9%	Samsung	1.6%
Intel	7.3%	Google	1.6%
unknown	6.9%	Oracle	1.5%
Novell	4.0%	Microsoft	1.4%
IBM	3.6%	AMD	1.3%
TI	3.4%	Freescale	1.3%
Broadcom	3.1%	Fujitsu	1.1%
consultants	2.2%	Atheros	1.1%
Nokia	1.8%	Wind River	1.0%

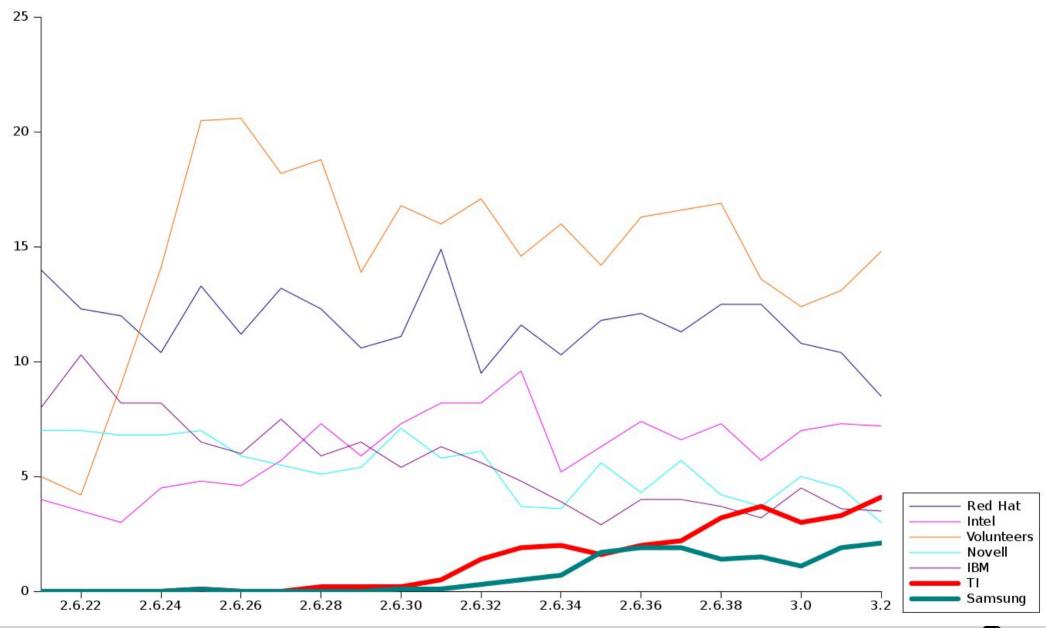


Kernel changeset contributions by employer





Kernel changeset contributions by employer





Also in February

Red Hat stops releasing individual kernel patches



March



2.6.38 released – March 14, 2011 (9,577 changes from 1198 developers)

Per-session group scheduling dcache scalability patch set Transmit packet steering Transparent huge pages Hierarchical block I/O bandwidth controller



Somebody needs to get a grip in the ARM community. I do want to do these merges, just to see how screwed up things are, but guys, this is just ridiculous. The pure amount of crazy churn is annoying in itself, but when I then get these "independent" pull requests from four different people, and they touch the same files, that indicates that something is wrong.

- Linus Torvalds, March 17



What is the "ARM problem"?

Wildly varying hardware "Embedded" mindset Little high-level oversight or communications

Results

Lots of little subtrees Lots of duplicated code A big ugly mess in general



Why is this happening

For years we have asked embedded vendors to contribute back to the kernel.



Why is this happening

For years we have asked embedded vendors to contribute back to the kernel.

...now they are doing it!



Cleaning up the mess

More high-level oversight Arnd Bergmann's arm-soc tree

More cleanup work
GPIO consolidation
Pinmux subsystem
Common clock framework

Move toward device tree

Eliminate lots of "board files"

Someday: one ARM kernel for all systems





Native Linux KVM Tool

A simple QEMU replacement Aimed at kernel developers

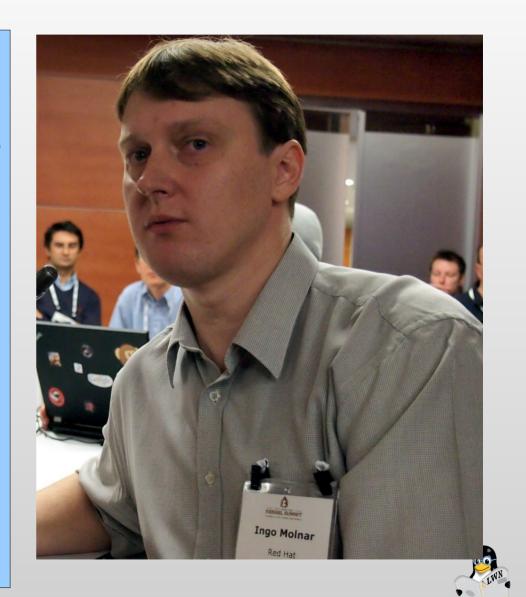
The sticking point:

The desire to add it to the kernel tree



I think it's only a matter of time until someone takes the Linux kernel, integrates klibc and a toolchain into it with some good initial userspace and goes wild with that concept, as a single, sane, 100% selfhosting and self-sufficient OSS project, tracking the release schedule of the Linux kernel.

Ingo Molnar, April 5



User-space code in the kernel tree?

Advantages
Wider visibility of the code

Develop ABI and users together

Encourage thinking across the boundary

Better integration



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Disadvantages

Kernel tree bloat

ABI stability problems

Other projects are disadvantaged

Where does it end?



April

The mobile space is about proprietary drivers

– Mark Charlebios, Qualcomm Innovation Center



May



Seccomp - sandboxing for Chrome A simple bitmask to limit available system calls

"Why not make it more powerful?"

Various filtering schemes proposed

Perhaps use tracepoints as enforcement points?

The end result Nothing merged



Yet another kernel release

2.6.39, May 18, 2011 (10,269 changesets, 1,258 developers)

Directed yield
IPset
Transcendent memory core
User namespace support
Media controller subsystem





During the 2.6.40 merge window

The voices in my head also tell me that the numbers are getting too big. I may just call the thing 2.8.0. And I almost guarantee that this PS is going to result in more discussion than the rest, but when the voices tell me to do things, I listen.

Linus Torvalds, May 23, 2011



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- Linus Torvalds, May 23, 2011

If you do this, I will buy you a bottle of whatever whiskey you want that I can get my hands on in Tokyo next week.

- Greg Kroah-Hartman



June



Ext4 snapshots posted

Save copies of a running ext4 filesystem

Useful for
System rollbacks
Backups
Factory reset

. . .



Why put all this effort into shoehorning in such a big an invasive feature to ext4 when btrfs does this all already? ...

The wonderful thing about ext4 is its a nice basic fs. If we're going to start doing lots of crazy things, why not do them to the fs that isn't yet in wide use and can afford to have crazy things done to it without screwing a bunch of users who already depend on ext4's stability?

- Josef Bacik



What's up with ext4?

"Bigalloc"

Allocate blocks in units >4096 bytes Makes operations on large files much faster Merged for 3.2

In the works

Snapshots
Inline data for small files
Secure erase support
Metadata checksumming

. . .



In other words

Ext4 will continue to develop and grow for a while yet.

I'm actually finding that ext4 has found a second life as a server file system in large cloud data centers. It turns out that if you don't need the fancy-shmancy features that Copy-on-Write file systems give you, they aren't free.

- Ted Ts'o



UEFI secure boot

The objective:

Only give control of the system to a "trusted" boot loader

This concept has value

Thwart bootloader rootkits

Ensure the system is running what you think it is

There is only one problem:



Who is "trusted"?

The owner of the computer?

The hardware vendor?

The software vendor?

The entertainment industry?



UEFI secure boot could easily be a mechanism by which we lose control of our computers.



Where things stand

Lots of work to call attention to the problem

Some concessions gained
All [x86] systems can be put into "setup mode"
It will be possible to install a signing key



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All [x86] systems can be put into "setup mode"
It will be possible to install a signing key

But:

Installing that key may not be easy
No provision for booting from CD
ARM systems can be totally locked down



July



3.0-rc7-rt

The first new realtime patch set since March Users had been stuck on 2.6.33



The state of realtime

Nice determinism on good hardware

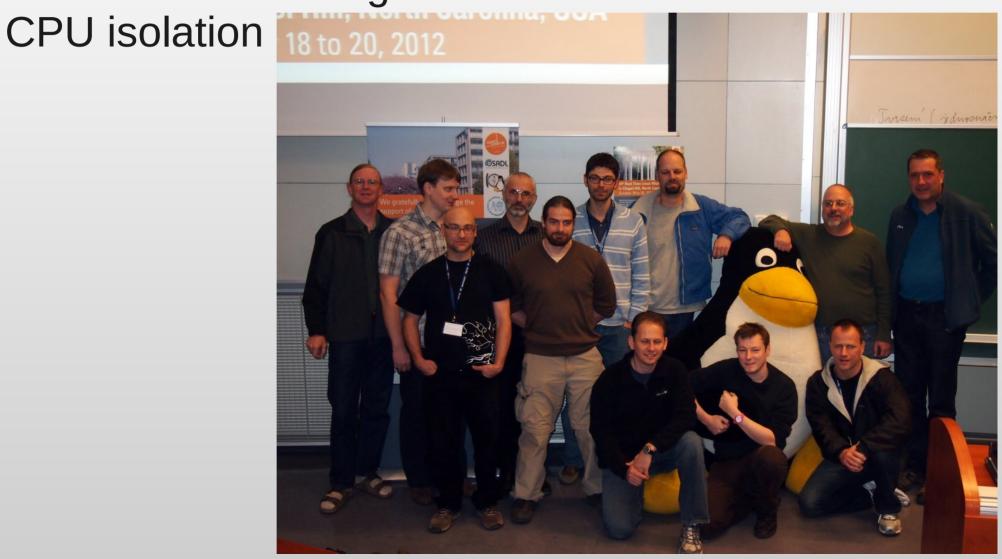
May have a solution on per-CPU data ...but involves scary locking assumptions

Plan is to merge most of it in the next year ...time will tell...



Open issues in realtime

Deadline scheduling



The 3.0 release is delayed

Nasty bug in the dcache scalability patches

The debugging crew:

Linus Torvalds

Al Viro

Hugh Dickins

...it still took them several days to figure it out



Some parts of the kernel have reached a truly scary level of complexity.



3.0 kernel released, July 21 (9,153 changes from 1,131 developers)

New POSIX clocks
BPF JIT compiler
sendmmsg() system call
ICMP sockets (unprivileged ping)
Namespace file descriptors
Cleancache







x32

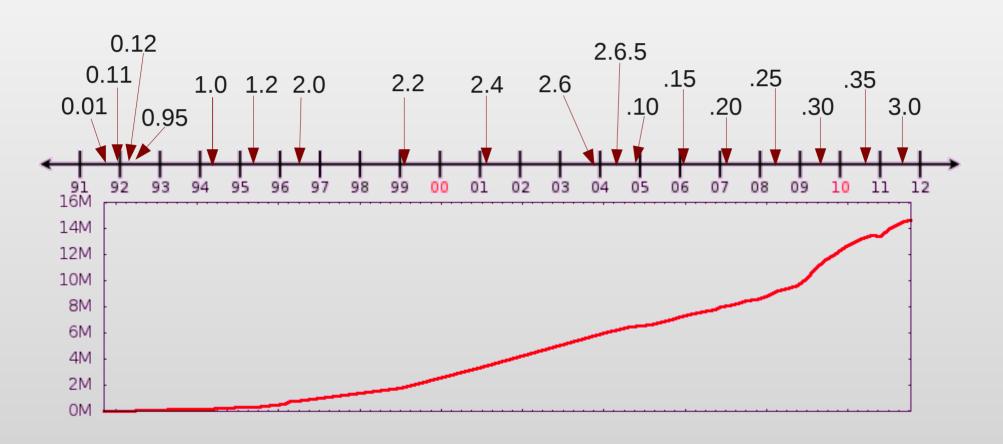
64-bit mode is great, but: 64-bit data is rarely needed Programs get larger, slower

The best of both worlds: the x32 ABI Run in full 64-bit mode
Use 32-bit data and pointers

Mostly a user-space problem But kernel support is needed



20 years of Linux





Kernel.org compromised

What is known:

Attackers had been on the system for some time Stolen credentials used; trojaned SSH installed Numerous associated machines compromised No attempts to corrupt software distribution

The immediate result:

kernel.org is down for almost two months The 3.1 kernel release is delayed



What has been done

A new kernel.org infrastructure has been built Lots of machines to separate functions New staff hired

Access has been restricted considerably "Maybe 450 shell accounts is a bad idea..."

A new kernel web of trust has been built

Vast support from the Linux Foundation



Still....

We do not take the security problem seriously enough.

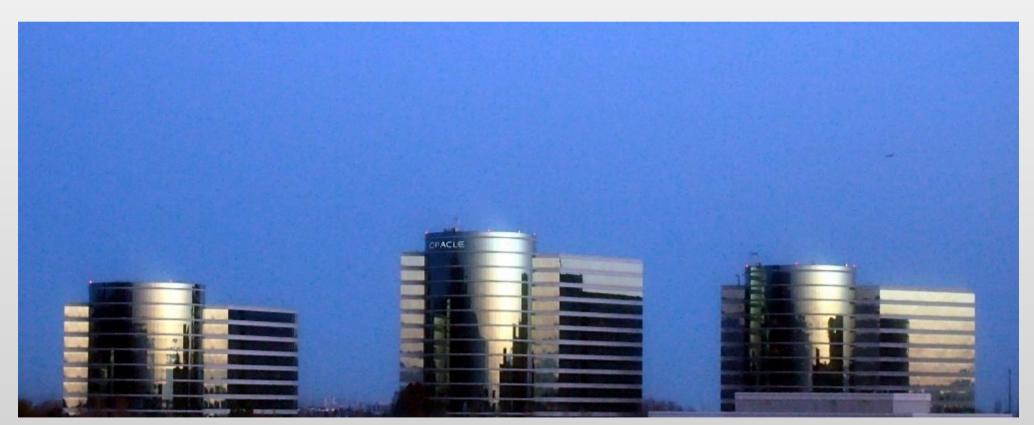


September



Oracle to use Btrfs by default

...sometime really soon now





Btrfs

Some new development work happening Lots of internal work Scrub feature

Stability is the biggest concern



Still missing

Btrfsck a hard problem, seemingly

Still under development.

Meanwhile
Root block history array
Read-only data recovery tool

Also missing: RAID 5/6 support Patches exist





2011 Kernel Summit





Two pivotal summit outcomes

- 1) Maintainers should say "no" more often
- 2) Widely-used code should be merged even if it is not up to normal technical standards

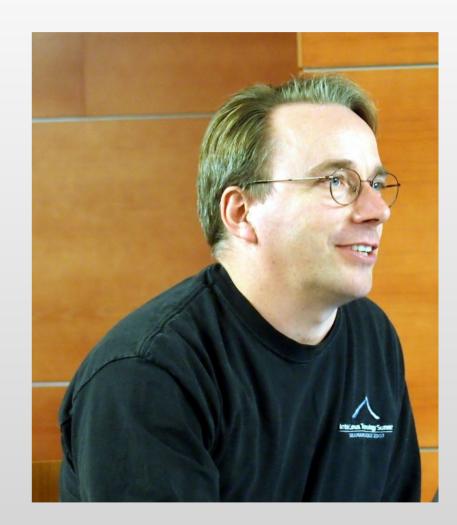


A slow moment at the Summit

The **3.1** kernel October 24, 2011 (8,693 changesets, 1,168 developers)

A 95 day cycle (average is 76)

Dynamic writeback throttling OpenRISC architecture PTRACE_SEIZE Iseek() hole finding



. . .



Embedded long-term support initiative



Two-year stable kernel maintenance
One kernel/year
Starting with 3.0

A separate tree for products Backports and such

A staging tree for upstreaming



November



Per-group TCP buffer limits

Limit kernel memory used by TCP buffers Accepted for 3.3

The first overt limit on kernel memory use Wanted for containers and such Lots more to come



Control groups

A simple mechanism for grouping processes ...that everybody hates

The real problem is the controllers

Memory usage (Now kernel memory usage too)

Block I/O bandwidth

Scheduling

CPU affinity

• • •

Expect a lot of cleanup work in this area



LTTng pulled into staging

A comprehensive tracing toolkit Widely used in some areas

Intended for merging into 3.3



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The outcome

LTTng loses





The Android mainlining project

An effort to get the Android kernel code merged

Includes

Binder - interprocess communication

Logger - user-space logging system

Low-memory killer

Pmem - contiguous memory allocation

RAM console

Timed GPIO

Ashmem - shared object storage



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January



Happy New Year

3.2 released, January 4, 2012 (11,828 changesets from 1,309 developers)

Proportional rate reduction
Extended verification module
CPU scheduler bandwidth controller
Cross-memory attach
Hexagon architecture
Btrfs recovery
I/O-less dirty throttling



3.3 merge window

"team" network device Network priority control group TCP buffer size controller Byte queue limits Open vSwitch

ARM LPAE support
The Android drivers return
DMA buffer sharing API

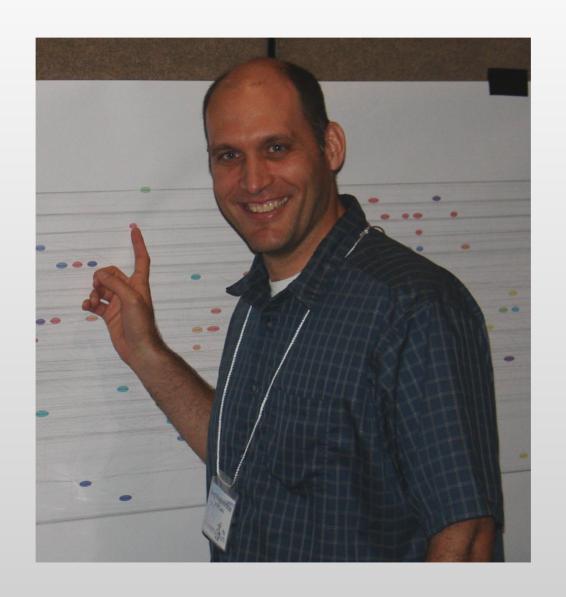
Expect 3.3 sometime in March



February



Greg KH joins the Linux Foundation





Stuff not covered

Writeback Transcendent memory Barriers Preemption disable Perf and ftrace Pin controller RAID x 4 Opportunistic suspend Power domains Common clocks Signed tags

Compaction stalls **GPL** violations **GPL** termination Patch review Testing tools SCSI targets **Bufferbloat** Power-aware sched. Solid-state storage IIO



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Perf and ftrace Ues Figh Sis? Pin controller RAID x 4 Opportunistic suspend Power domains Common clocks

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