# Controlling Memory Fragmentation and Higher Order Allocation Failure: Analysis, Observations and Results



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# <u>INTRODUCTION</u>

What is Memory Fragmentation?

When a Linux device has been running continuously over a time without reboot and keeps allocating and de-allocating pages, the pages become fragmented. The bigger contiguous free pages become zero and free pages are only available in many smaller pages which are not contiguous. Thus even if we have lots of free memory in smaller units, the page allocation in kernel may fail.

This typical problem is called "External Memory Fragmentation" here after referred to as memory fragmentation.

# **INTRODUCTION**

- Effect of Memory Fragmentation :
  - ➤ Memory Fragmentation can cause a system to lose its ability to launch new process.
  - ➤ Memory fragmentation becomes more of an issue in embedded devices and Linux mobiles.
    - DRAM + Flash , Swapless system
  - Memory fragmentation can become more critical with high multimedia and graphics activities which requires contiguous higher-order pages.

# MEASURING FRAGMENTATION LEVEL

- It is important to measure fragmentation level across each zones and for each higher-order allocation in kernel.
- We believe by measuring fragmentation level during page allocation we can control higherorder allocation failure in kernel.
- We developed kernel utility to measure fragmentation level during runtime without enabling memory COMPACTION.

# MEASURING FRAGMENTATION LEVEL

Formula to measure fragmentation level in percentage :

$$FragLevel(\%) = \frac{TotalFreePages - \sum_{i=j}^{N} (2^{i}.k_{i})}{TotalFreePages} X100$$

```
TotalFreePages = Total number of free pages in each Node
```

N = MAX\_ORDER - 1 → The highest order of allocation

j = the desired order requested

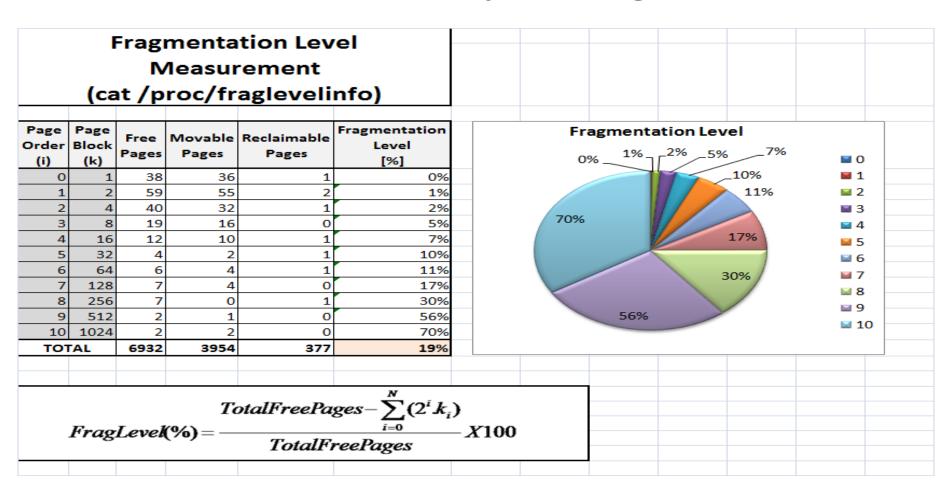
i = page order → 0 to N

Ki = Number of free pages in ith order block

( The above formula derived from Mel Gorman's paper : "Measuring the Impact of the Linux Memory Manager")

# MEASURING FRAGMENTATION LEVEL

# SAMPLE OUTPUT: cat /proc/fraglevelinfo



## **Example:**

### Lets say a NORMAL zone looks like this at some point of time

		<b>2</b> °	2 <sup>1</sup>	2 <sup>2</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>7</sup>	<b>2</b> <sup>8</sup>	<b>2</b> <sup>9</sup>	<b>2</b> <sup>10</sup>
Node 0, zone	Normal	3	20	104	13	13	1	1	1	0	0	0

Now lets apply our formula to measure fragmentation level for order 2<sup>5</sup>.

Here , TotalFreePages = 
$$(3x1 + 20x2 + 104x4 + 13x8 + 13x16 + 1x32 + 1x64 + 1x127) = 994$$

### Therefore;

```
% Fragmentation = (994 - [(2^5)*1 + (2^6)*1 + (2^7)*1 + (2^8)*0 + (2^9)*0 + (2^10)*0])*100)/994
```

%Fragmentation = (770 \* 100) / 994 = 77.46 % → 77 % (round off)

# MEMORY FRAGMENTATION ANALYSIS

- We developed a sample kernel module and test utility to perform higher-order allocation and doing memory fragmentation analysis before and after the allocation.
- Test utility developed and tested for Samsung Linux Platform (kernel2.6.32 and kernel2.6.36)
- Test Utility can be shared on Linux Test Project (LTP) after further improvements.

### **JUST AFTER PHONE BOOT-UP (Samsung Linux Kernel 2.6.36)**

```
/opt/home/root # cat /proc/fraglevelinfo
Node:0, Zone:Normal
order
         FreePages MovablePages ReclaimablePages Fragmentation[%]
                                                                      0%
                                                                      0%
  1
                             1
  2
                             1
                                                                      0%
                                                   2
  3
4
5
6
                             1
                                                                      0%
                                                                      0%
                                                                      0%
                             1
                                                                      0%
  7
                             1
                                                                      0%
  8
                             O
                                                                      1%
                                                                      1%
  9
               1
                                                                      1%
10
             102
                             101
TotalFreePages: 106414
TotalMovablePages: 103678
TotalReclaimablePages: 731
Overall Fragmentation: 0%
Node:0, Zone:HighMem
         FreePages MovablePages ReclaimablePages Fragmentation[%]
order
  O
              15
                                                                      0%
  1
              15
                                                                      0%
                                                   O
  2
3
                                                                      0%
              11
                                                   O
                                                                      0%
              10
  4
5
6
              10
                                                   O
                                                                      0%
                                                                      0%
              11
                                                   О
                                                                      1%
  7
                                                                      1%
                                                   O
  8
                             1
               1
                                                                      3%
                                                   0
  9
               1
                                                   O
                                                                      3%
10
                                                   O
                                                                      4%
              55
TotalFreePages: 59049
TotalMovablePages: 56665
TotalReclaimablePages: 0
Overall Fragmentation: 1%
```

/opt/home/rod	t # cat /pr	roc/budo	dyinfo										1
Node 0, zone	Normal	2	2	5	3	4	2	6	5	1	1	102	
Node 0, zone	HighMem	15	15	11	10	10	11	4	8	1	1	55	

# AFTER RUNNING VARIOUS APPLICATIONS (Browser, WiFi Video Share, Camera, Voice Recorder, eBooks, Few Games) for ½ an Hour and then killing All)

/opt/pi	ntu # cat /	proc/fraglevel	info	
Node:0.	Zone:Norma	1		
order	FreePages	MovablePages	ReclaimablePages	Fragmentation[%]
0	1139	800	1	0%
ĭ	795	540	7	1%
5	500	345		3%
5	315	229	1	5%
2 3 4 5 6 7	237	168	<u> </u>	8%
4		70	1	
2	114		1	13%
9	67	26	<u> </u>	17%
	25	15	1	23%
8	35	10	2 1 0 1 1 1 0	26%
9	16	9		37%
10	42	36	0	47%
TotalFr	eePages: 82	333		
TotalMo	vabléPages:	57636		
TotalRe	claimab <b>í</b> ePa	iges: 511		
overal1	Fragmentat	ion/ 16%		
Node:0.	Zone:HighM	lem		
order	FreePages	MovablePages	ReclaimablePages	Fragmentation[%]
0	2281	1994	0	0%
	1382	1140	Ō	9%
2	861	781	Ö	21%
3	474	430	ŏ	36%
Ā	226	193	ŏ	52%
5	106	85	ŏ	67%
6	38	26	ŏ	82%
1 2 3 4 5 6 7	8	6	ŏ	92%
8	3	ŏ	ŏ	96%
9	ŏ	ŏ	ŏ	
	ŏ	ŏ		100%
10		_	0	100%
TotalFr	eePages: 23	513		
TOTAIMO	vabļēPages:	190/8		
TotalRe	claimab1ePa	iges: 0		
overall	Fragmentat	iion 59%		

/opt/pintu # Node 0, zone	cat /proc	/buddyi	nfo									42 0
Node 0, zone	Normal	1139	795	500	315	237	114	67	25	35	16	42
Node 0, zone					474	226	106	38	8	3	0	0
												11

### **ALLOCATION SUCCESS CASE - (Samsung Kernel 2.6.32)**

/opt/pintu # ls -l /dev/pinchar

crw----- 1 root root 10, 49 Jan 12 13:21 /dev/pinchar

### /opt/pintu # ./app\_pinchar.bin

		<b>2</b> °	2 <sup>1</sup>	2 <sup>2</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>7</sup>	<b>2</b> <sup>8</sup>	<b>2</b> <sup>9</sup>	2 <sup>10</sup>
Node 0, zone	DMA	3	1	1	1	0	0	0	1	0	0	0
Node 1, zone	DMA	197	129	0	0	0	0	0	0	0	0	0
Node 2, zone	DMA	30	12	14	4	5	8	6	4	2	1	27

Enter the page order(in power of 2):	16	2^4 order block	16 x 4K = 64K bytes
Enter the number of such block :	5		

### State After The Allocation Request Is Successful

		<b>2</b> °	<b>2</b> <sup>1</sup>	2 <sup>2</sup>	2 <sup>3</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>7</sup>	<b>2</b> <sup>8</sup>	<b>2</b> <sup>9</sup>	2 <sup>10</sup>
Node 0, zone	DMA	3	1	1	1	0	0	0	1	0	0	0
Node 1, zone	DMA	169	143	1	0	0	0	0	0	0	0	0
Node 2, zone	DMA	19	12	12	5	2	7	6	4	2	1	27

### **Explanation:**

5 of 2^4 order blocks were requested. These request could only be satisfied by Node 2, Thus Node 2 were selected for allocation. But in Node 2 also, out of 5 (2^4) blocks only 3 could be allocated. Then the other 2 were allocated by splitting the 2^5 order blocks.

$$5 \times 2^4 = [3 \times 2^4 + (1 \times 2^5)] = [3 \times 2^4 + (1 \times 2 \times 2^4)] = [5 \times 2^4]$$

### **ALLOCATION FAILURE CASE - (Samsung Kernel 2.6.32)**

### **Buddy State Before Allocation**

		<b>2</b> °	<b>2</b> <sup>1</sup>	2 <sup>2</sup>	2 <sup>3</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>7</sup>	<b>2</b> <sup>8</sup>	<b>2</b> <sup>9</sup>	2 <sup>10</sup>
Node 0, zone	DMA	1	0	0	0	1	0	2	0	0	0	0
Node 1, zone	DMA	247	181	15	1	1	2	0	9	10	8	11
Node 2, zone	DMA	19	19	11	3	6	3	4	2	2	2	29

Enter the page order(in power of 2):	1024	2^10 order block	1024 x 4K = 4096K bytes	
Enter the number of such block:	50	(thi	is is the highest order)	

ERROR: ioctl - PINCHAR\_ALLOC - Failed, after block num = 48!!!

**Explanation:** As you can see the allocation request of 1024 x 50 pages is failed after 47 such allocation. But still there were enough free pages available in lower order.

### **Buddy State After Allocation FAILED And Other Allocation Freed**

		<b>2</b> °	2 <sup>1</sup>	2 <sup>2</sup>	<b>2</b> <sup>3</sup>	<b>2</b> <sup>4</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>7</sup>	<b>2</b> <sup>8</sup>	<b>2</b> <sup>9</sup>	2 <sup>10</sup>
Node 0, zone	DMA	1	0	18	9	0	0	0	0	0	0	0
Node 1, zone	DMA	88	77	36	25	25	43	23	19	18	11	20
Node 2, zone	DMA	18	14	10	4	5	2	3	3	2	2	29

**Explanation:** The interesting think to note here is that after requested allocation was failed, kernel tried to arrange that many blocks in desired order so that next similar request can be succeeded.

# **Observations**

- \_\_alloc\_pages\_nodemask : This is the heart of all allocation in kernel.
- We measure fragmentation level for each higher order here.
- Track higher-order allocation during high fragmentation. Anything above PAGE\_ALLOC\_COSTLY\_ORDER(==3) is considered higher-order allocation and becomes critical.

# **Observations**

- Direct reclaim does some progress but still could not return any pages during first run.
- Similarly direct compact (from kernel2.6.36 onwards) is helpful but still not effective for very higher-order allocation.
- May be the other way round (first direct\_reclaim then direct\_compact) could be more helpful.
- A small amount of delay(for GFP\_KERNEL) is required after direct\_{reclaim,compact} and before retry.
   Maybe due to lazy buddy allocator.

# **Experiments Results**

- We performed some experiments with higher-order allocation and got some results.
- We found that whenever we run any application "Xorg" perform 4 or 8 order allocation.
- The browser always requires order-4 allocation.

```
/opt/pintu # ps ax | grep browser
7159? Ssl 0:03 /opt/apps/com.samsung.browser/bin/browser

[ 3830.215613] [HIGHERORDER_DEBUG] : __alloc_pages_nodemask is called by process <PID = 1168, NAME = Xorg> !!!
  [ 3830.227243] [HIGHERORDER_DEBUG] : ZONE : Normal, NODE : 0, ORDER = 8, Fragmentation Level = 29%
  [ 3830.235645] [HIGHERORDER_DEBUG] : __alloc_pages_nodemask is called by process <PID = 1168, NAME = Xorg> !!!
  [ 3830.244575] [HIGHERORDER_DEBUG] : ZONE : Normal, NODE : 0, ORDER = 4, Fragmentation Level = 13%
  (Around 10 times)

[ 3831.355884] [HIGHERORDER_DEBUG] : __alloc_pages_nodemask is called by process <PID = 7159, NAME = browser> !!!
  [ 3831.364649] [HIGHERORDER_DEBUG] : ZONE : Normal, NODE : 0, ORDER = 4, Fragmentation Level = 13%
  [ 3831.373484] [HIGHERORDER_DEBUG] : __alloc_pages_nodemask is called by process <PID = 7159, NAME = browser> !!!
  [ 3831.383134] [HIGHERORDER_DEBUG] : ZONE : Normal, NODE : 0, ORDER = 4, Fragmentation Level = 13%
  (Around 26 times)
```

### RESULT #1 - (Samsung Kernel 2.6.32)

[24770.686688] < PINCHAR > : PINCHAR ALLOCATE - Success(index = 0)!

		<b>2</b> °	2 <sup>1</sup>	2 <sup>2</sup>	<b>2</b> <sup>3</sup>	2 <sup>4</sup>	<b>2</b> <sup>5</sup>	<b>2</b> <sup>6</sup>	<b>2</b> <sup>7</sup>	2 <sup>8</sup>	<b>2</b> <sup>9</sup>	2 <sup>10</sup>
Node 0, zone	DMA	685	104	1	1	0	0	0	1	0	0	0
Node 1, zone	DMA	33	19	31	18	9	1	1	0	0	0	0
Node 2, zone	DMA	11	50	9	5	5	3	2	1	0	0	0

Enter the page order(in power of 2):	1024	2^10 order block	1024 x 4K = 4096K bytes
Enter the number of such block:	10	(thi	is is the highest order)

**Explanation:** As you can see here, due to 100% fragmentation, page allocation request was failing, even after direct reclaim (slow path). But after a delay and retrying allocation request again, all subsequent allocation were successful.

This delay indicates something needs to be done after direct reclaim. Maybe wait till lazy buddy allocator arranges free pages in the subsequent free areas.

### **RESULT #2 - (Samsung Kernel 2.6.36) [Without COMPACTION]**

### **Initial Fragmentation Level**

/opt/pint	u # cat /proc	/fraglev	/elinfo ;	cat /	proc/bu	ddyinfo						
order 0 1 2 3 4 5 6 7 8 9 10 TotalFree TotalMova TotalRecl	one:Normal FreePages 271 171 143 100 154 79 26 7 10 0 1 Pages: 13117 blePages: 575 aimablePages:	2756	lePages 5 5 1 1 2 1 3 0 0	Rec	laimabl 0 0 1 18 61 27 6 3 0 0	ePages			ation[% 0% 2% 4% 9% 15% 33% 53% 65% 72% 92%	6]		
Order  0 1 2 3 4 5 6 7 8 9 10 TotalFree TotalMova TotalRecl	one:HighMem FreePages 6850 8826 261 9 5 0 0 0 0 Pages: 25698 blePages: 2555 aimablePages:	683 880 25	)3	Rec	laimabl 0 0 0 0 0 0 0 0	ePages		1 1 1 1 1	ation[% 0% 26% 95% 99% 00% 00% 00% 00% 00%	6]		
Node 0, z Node 0, z	one Normal one HighMem	271 6850	171 8826	143 261	100 9	154 5	79 0	26 0	7 0	10 0	0	1

### **After Higher order allocation request**

#./app\_pinchar.bin 1024 25

```
[17949.789934] [HIGHERORDER DEBUG]: alloc pages nodemask is called by process <PID = 27713, NAME = app pinchar.bin>!!!
[17949.801633] [HIGHERORDER DEBUG] : ZONE : Normal, NODE : 0, ORDER = 10, Fragmentation Level = 92%
[17949.811073] [HIGHERORDER DEBUG]: alloc pages nodemask: Allocation going via - slowpath!!!
[17949.831793] [HIGHERORDER DEBUG] : did some progress = 151
[17949.844090] [HIGHERORDER DEBUG] : NO pages......even after direct reclaim
[17949.859104] app pinchar.bin: page allocation failure. order:10, mode:0x40d0
[17951.879156] [HIGHERORDER DEBUG] : Trying - Final time !!!!!!!!!
[17951.893248] <PINCHAR> : PINCHAR ALLOCATE - Success(index = 0)
[17960.189583] [HIGHERORDER DEBUG]: alloc pages nodemask is called by process <PID = 27713, NAME = app pinchar.bin>!!!
[17960.201128] [HIGHERORDER DEBUG] : ZONE : Normal, NODE : 0, ORDER = 10, Fragmentation Level = 98%
[17960.210269] [HIGHERORDER DEBUG]: alloc pages nodemask: Allocation going via - slowpath!!!
[17960.335044] [HIGHERORDER DEBUG] : did some progress = 887
[17960.339918] [HIGHERORDER DEBUG] : Got some pages after direct reclaim .....
[17960.368939] < PINCHAR > : PINCHAR ALLOCATE - Success(index = 4)!
[17964.518845] [HIGHERORDER DEBUG]: alloc pages nodemask is called by process <PID = 27713, NAME = app pinchar.bin>!!!
[17964.530629] [HIGHERORDER DEBUG] : ZONE : Normal, NODE : 0, ORDER = 10, Fragmentation Level = 83%
[17964.547138] < PINCHAR > : PINCHAR ALLOCATE - Success (index = 8)!
[17965.552976] [HIGHERORDER DEBUG]: alloc pages nodemask is called by process <PID = 27713, NAME = app pinchar.bin>!!!
[17965.564319] [HIGHERORDER DEBUG] : ZONE : Normal, NODE : 0, ORDER = 10, Fragmentation Level = 84%
[17965.580823] < PINCHAR > : PINCHAR ALLOCATE - Success(index = 9)!
[17966.586440] [HIGHERORDER DEBUG]: alloc pages nodemask is called by process <PID = 27713, NAME = app pinchar.bin>!!!
[17966.597175] [HIGHERORDER DEBUG]: ZONE: Normal, NODE: 0, ORDER = 10, Fragmentation Level = 85%
[17966.613424] < PINCHAR > : PINCHAR ALLOCATE - Success(index = 10)!
```

Allocation failed directly during the first attempt itself even after direct reclaim. But after introducing a delay and retrying, all further allocation succeeded. May be Kswapd takes sometime to clear up dirty pages and buddy adding it back to free area.

### **Final Fragmentation Level**

Node:0 Order 0 1 2 3 4 5 6 7 8 9 10 Totalk Totalk	Free Free 86 86 86 56 35 22 9 7 1 2 reePage lovableP	Pages 4 3 8 9 6 8 2 0 8	Movak 58 62 49 43 28 18 6 1	olePages 55 33 26 93 34		/proc/bu 135 118 107 74 41 28 16 6 0	-		Fragmen	tation[9 0% 0% 2% 5% 10% 18% 28% 41% 51% 69% 75%	6]		
Order 0 1 2 3 4 5 6 7 8 9 10 TotalF	457 823 478 176 36 3 reePage lovableP	Pages 9 9 9 0 2	452 820 475 173 35 83 0	)2 56 34	Red	claimabl 0 0 0 0 0 0 0 0 0	ePages		:	tation[9 0% 7% 34% 65% 88% 97% 99% 100% 100%	6]		
	, zone , zone	Normal HighMem	864 4579	862 8239	868 4789	659 1760	566 362	358 35	222 4	90 0	78 0	13 0	26 0

Here you can see lots of movable pages after lots of direct reclaim. Thus direct compact might be helpful after direct reclaim and not before.

### **EXPERIMENTATION DATA**

Page Order	Block Used	Available Blocks	No of Blocks Requested	Current Fragmentation Level	No of Blocks Allocated	Pass Rate
10	1024	0	20	100%	20	100%
9	512	11	20	94%	20	100%
8	256	4	20	90%	20	100%
8	256	0	50	100%	50	100%
9	512	1	30	97%	30	100%
10	1024	28	40	10%	40	100%
10	1024	0	50	100%	46	92%

### DATA COLLECTED ON:

Samsung Mobile Target With Kernel 2.6.32

# **SUMMARY**

- Measuring fragmentation level and tracking higherorder is important at least for low memory notifier.
- It was observed that allocation takes slowpath whenever fragmentation level is above 90%.
- The delay introduced here is only for experimental purpose.
  - > Delay could be because, dirty pages has to be written to the disk before it is marked freed.
  - ➤ May be the real thing could be to wait till lazy buddy allocator rearranges the free pages.
  - > This is valid only for GFP\_KERNEL where a sleep is allowed.

- For fragmentation > 90%, introduce temporary kernel thread to do direct reclaim/compact in background.
  - > Buy the time you come back for next request, pages will be ready for you.
  - Not enough data to share. Further experiments in progress.
- From kernel2.6.35 COMPACTION contains its own fragmentation level measurement.
  - /sys/kernel/debug/extfrag/unusable\_index
  - > But this requires COMPACTION and HUGETLB to be enabled.
  - May be we can utilize this from kernel2.6.35 onwards.
  - > Difficult to back port compaction to lower kernel version.
  - Mostly helpful for large system and may not be useful for small embedded products.

- Can we introduce something like system wide fragmentation level ???
- Reboot is not a good choice for end users even for small system.
- May be introduce something like "Reset Physical memory state".
  - Bring back memory to original state without reboot.
  - Not enough data.
  - ➤ May be develop system utility to shrink physical memory using "shrink\_all\_memory" used during snapshot image creation in Hibernation.

- Reserving memory during boot time can reduce fragmentation to some extent.
  - Good only if you have bigger RAM.
  - > Tracking higher-order fragmentation level can help decide which memory to reserve in future.
  - May be something like dynamic reserving based on past performance could be better.
- Contiguous Memory Allocator (CMA) and Virtual Contiguous Memory (VCM) can help fight fragmentation.
  - CMA: same like reserving memory during boot but transparently allows the memory to be reused and latter migrate pages to create similar chunk.
  - Can be used for frame buffers and other memory hungry multimedia devices.

- But CMA requires pages to be movable and may now use compaction, again not guaranteed because most kernel pages are not movable.
  - May be we have to limit the reuse of CMA region.
  - Or share CMA region only for very high order.
- Problem easily reproducible in DRAM + Flash swapless embedded system without HighMem.
- Further combination of investigation is in progress to derive a solution which does not requires reboot.

# **Some References**

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# Thank You !!!