

# An Electronic Digital Computor

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## An Electronic Digital Computor

Using Cold Cathode Counting Tubes for Storage

By R. C. M. Barnes, B.Sc., E. H. Cooke-Yarborough, M.A., A.M.I.E.E., and D. G. A. Thomas, M.A. \*

A small sequence-controlled computer is described; it has storage capacity for up to 90 numbers cach of eight electional digits. Cold catabode counting tubes (Debatrons) ore used as storage elements. Parallel operation is employed, and the speed of operation is comparable with that of a desk calculating machine. Apart from speed, however, the computer provides most of the facilities found in larger, faster digital computors. Information is fed into the machine from perforated paper tages and at elegizative or rate perforator records the results.

SPERAL large-sele computers have been described in Seventy states, making the ENMAC"—with shard where circuits, the EDSAC"—with mercury delay storage, the Underveity of Manchester computers'—with cathode-ray storage and the Birkbeck College computer with magnitive carry out eneromes computations, too lengthy to perform by other meant. In addition smaller scale relay computers'—have been built for general with. The medium proports' have been built for general with. The medium control of the scale relay to the state of the scale relation of the scale relation to the scale of the scale of the lations for which the large computers have been designed. It is intended rather to do the work of a few operations with

place. In the present machine these two parts are almost entirely separate. Because a relatively low operating speed is accepted, it

Because a relatively low operating speed is accepted. It becomes practical to read orders and other data from the becomes practical to read orders and other data from the stone in the BTL machine <sup>1-1</sup>3. Design is thus simplified and much less stone in the BTL machine <sup>1-1</sup>3. Design is thus simplified and much less storage capacity is needed. The low operating speed also makes possible the use of relays for routing and sequence control. Since each relay may carry many contacts, this leads to simplicity in setting up connection, reproved reliability is of great importance.

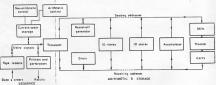


Fig. 1. Block schematic diagram of computer.

desk calculating machines where the work to be done is of a routine or repetitive nature. As a result, simplicity and reliability have been regarded as of more importance than speed of operation and the machine performs the operations of addition, subtraction, multiplication and division little faster than a desk calculator.

little faster than a desk calculator.
Fig. 1 is a sphematic diagram of this machine. The
design can, for convenience, be regarded as consisting of
two parts. The first includes the means by which the
machine receives orders, controls the sequence of operaction of the control of the sequence of operaction of the control of the control of the control
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Numerical data obtained in the course of the calculation motive design in a further as may be considered, largely would be reacted. Other methods were considered, largely for their needed. Other methods were considered, largely for their needed of the considered o

<sup>\*</sup> Electronics Division, A.E.R.E., Harvell.

sources is much reduced by comparison with other storage methods. Use of these tubes permits the computor to work on a decimal, rather than a binary, basis. This is convenient from several points of view; it is particularly advantageous in a parallel computor, since the number of digits and therefore the number of parallel arithmetical

units, is considerably reduced.

Provided that the number of stores required is not too large, this method of storage is thought to compare favourably with other methods on the basis of cost, size, and power consumption per digit. When the computor is fitted with its full storage capacity (90 eight-digit numbers, equivalent to 2,700 binary digits) just over 80% nor 100 binary digits) just over 80% nor 100 binary digits) just over 80% nor 100 binary digits in the 100 binary digits in 100

current and from life tests to date, tube life is expected to be many thousands of hours. Each storage address consists of nine tubes, corresponding to eight digits with one additional tube to indicate the

sign. The decimal point is regarded as being after the first figure of the number stored, whether this is a significant figure or zero, and the capacity n of each store is + 10 > n > -10. In the capacity n of each store is + 10 > n > -10. In the capacity n of each store is + 10 > n > -10. In the capacity n of each store is + 10 > n > -10. In the capacity n of each store is + 10 > n > -10. In the capacity n of each store is + 10 > n > -10. In the capacity n of each store is a store of the capacity of each store is the capacity of each store is the capacity of each store is each store in the capacity of each store is each store in the capacity of each store is each store in the capacity of each store is each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacity of each store is each store in the capacity of each store in the capacit

handled as compenents except on the punched tape and when finally punched tape and when finally and sign. The basic process of addition for subtraction by addition of the complements) is carried out electronically in order to save the wear that would occur if high speal relavision is by a sequence of multiple additions for subtractions), with the poperation of a relay circuit to shift the decimal point as each multiple addition takes pilice.

off tables of functions and the low operating speeds will often make it impracticable to calculate the required functions ab initio. This must be regarded as a limitation of the present machine; it may eventually be overcome by providing a special table-reading attachment.

reading attachment.
The computor has been operating with 20 eight-digit stores for some months. This has made possible the establishment of most of the electronic and relay circuits. Although the machine has not yet been put into service it seems appropriate at this time to publish details of the design

which has been established.

The computor contains about 180
relays, 18 Dekatrons, 80 thermionic valves and 40 cold
cathode triodes, plus 28 relays and 90 Dekatrons per
ten stores. It occupies three 7ft, Post Office racks, together
with an additional rack per four groups of ten stores, one
smaller rack for the power supplies and a table for tape
readers and printers. Total power occupantion is about 1kW.

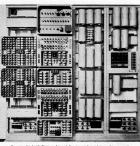
#### Facilities

The computor operates solely under the control of orders and data fed in from perforated paper tapes, and no manual switching or plugging up is necessary. Provision is made for a maximum of eight input tapes read by separate tape readers. Each tape may contain orders, input data, and constants. A typical allocation of tapes is expected to be a master programme tape, a tape giving the data relevant to the particular problem and a number of closed loops containing the longer sub-routines.

The output of the computor is fed into either a perforator or a modified teleprinter operating as a page printer. A maximum of eight printers and perforators can be accommodated although only a few of the positions are used at present. It is possible to feed the perforated tape from an output perforator into an input tape reader to give a long-term storage when the nature of the work.

permits.

The first order is always read from tape reader No. 01 and subsequent orders come from the same source until an order specifies a change. If the source of orders is changed to a Dekatron number store, orders are drawn from successive stores until the sequence is interrupted. Orders are normally only stored in the Dekatron stores when they take the form of short sub-routines or if they are to be



Computer fitted with 20 stores, but excluding tape readers, printers and power supply. The left hand rack carries two units each of ten stores. The left centre rack carries the electronic pulse generator (top), the arishmetic units with relay shift unit (centre) and the translator unit (bottom). The two right hand racks contain the relay sequence.

subjected to arithmetical operations; the majority of orders will be read directly from the tape readers.

Orders appear on the tape as blocks of five (decimal) digits and each digit is represented by one row of perforations. Numerical data are perforated as blocks of eight digits preceded by a sign character. Blocks of digits are separated by a character which can be either a space or one of ten block-marking characters. A search facility is one of the block-marking characters. A search facility is along its hap until it finish the required block-marking character.

ARITHMETICAL ORDERS
Arithmetical orders consist of one digit defining the

arithmetical operation to be performed, a sending address (two digits) and a receiving address (two digits). The sending address may be a Dekatron store, the accumulator, a tape reader, or the rounding-off circuit. The receiving address may be a store, the accumulator, a printer or perforator, or a drain for unwanted numbers. Seven operation digits have been allocated as follows:

 Add the contents of the sending address into the receiving address, and leave the contents of the sending address unchanged.

Add the contents of the sending address into the receiving address and clear the sending address.
 Subtract and do not clear the sending address.
 Subtract and clear the sending address.

(4) Subfract and clear five sending address by the (5) Multiply the contents of the sending address by the contents of the "receiving address aclearing the contents of the "receiving address, containing," (6) Divide the contents of the accumulator by the contents of the sending address, form the quotient in the receiving address and leave the remainder in the

(7) Add the positive modulus of the contents of the sending address into the receiving address.

The southing and receiving addresses must not lies in the same group of the store; in the case of multiplication and division they must not refer to an input or output organ. It should be noticed that addition and substraction do not require the use of an accumulator, because the result of a division of the store of

NON-MATTHMETICAL ORDERS
In non-arithmetical orders the operation digit is "0" and is followed by two further digits calling for a change of the address from which orders are obtained, search for a block-marking character, examination of the sign of the number in a specified address, or selection of the layout of the printed result. The last two digits specify the address.

concerned.

The sign examination facility allows the future action of
the computor to be controlled by the state reached by the
calculation. Any store can be examined for the presence
of either a positive or a negative sign. If the sign looked
for is found, an affirmative relay is, locked in and any

subsequent conditional orders (see below) are performed; otherwise they are ignored.

The order calling for a change in the order source can be made conditional or unconditional; if conditional it is only performed if the previous sign examination has given an

made conditional or unconditional; if conditional it is only performed if the previous sign examination has given an affirmative result. This allows a new set of orders to be brought in when the previous set has brought the calculations to a desired stage.

The order requiring a search for a block-marking behavior of the property of the property of the property of the line between the present position to the ignored which line between the present position that the constitution of the result of the previous sign examination. Facilities are provided for the tage to be searched in either the forward available only allow search in the forward direction. The ability to search in reverse will allow sections of the order to the present position of the present property of the present property of the present property of the present present property of the present present present present present present present present the present present present present present present present present the present present present present present present present present the present present present present present present present the present present present present present present present present the present present present present present present present present the present present present present present present present the present present present present present present present present the present present present present present present present present the present present present present present present present present the present present present present present present present present the present prese

The print layout is variable to a limited extent. There are ten layouts available for printing each eight digit number, for instance one such layout might consist of the eight digits followed by spaces, and another the eight digits followed by spaces, and another the eight digits followed by inse-feed and carriage return. An order is required to define which layout is to be used, and it stays in force until realseed by another print layout order.

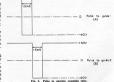
#### Arithmetical Operations

ADDITION AND SURTRACTION WITH STROLLE TURES
As has already been stated, addition is the fundamental
operation. Subtraction is effected by adding the complement, and multiplication (or division) by a sequence of
multiple additions (or subtractions). The circuits used for
adding a number from one store into another therefore
form the basis of the arithmetic unit, and are described in
detail. Numerical information is handled in the form of



trains of up to nise pulses. Each train represents one doctomal digit, and a spraited operation is employed there is a separate channel to handle each digit. For normal arrathrectical operations the pulses are routed, through a ratherated and persons the pulses are routed, through a real content of the results of the real content of the real co

It is first necessary to consider the mode of operation of the gas-filled tubes (Dekatrons) which form the basis



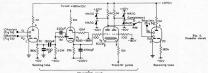
of the storage system. The arrangement of a Dekatron is shown diagrammatically in Fig. 2. Ten cathods are equally disposed round a central anode, and all but one are connected together. Between adjacent cathodes there are two guide electrodes, and these are connected together in the way shown.

Normally the zero cathode is taken through a resistor to earth or to a small bias potential, and the remaining cathodes are earthed. The guides rest at +60V and the anode is connected through a resistance of  $1 \, \mathrm{M}\Omega$  to +370V.

Pales of the form shown in Fig. 3 are applied to the guide electrodes to move the discharge from one cathode to the next. During the first palse the discharge moves to the next pales of the discharge moves on the discharge moves on to guide 2, which is them more negative than the previous cathode. At the end of the second pulse, advanced one step. It is not necessary for the two guide pulses to overlap, but it is important that any time interval advanced one step. It is not necessary for the two guide pulses to overlap, but it is important that any time interval second should not be comparable with the de-ionization right anode the complement on 10. These pulses are sustable for direct connexion to one guide of the receiving tube. Since, for the purpose of subtraction, the complement on 9 is required, it is arranged that only 9 pulses are applied to the other guide on the receiving tube (Fig. 56g) so that the first pulse of the train shown in Fig. 3(e) is ineffective in moving the discharge.

in moving the discharge.

If the sending tube is to be cleared to zero during the transfer, the complement output of the transfer unit is connected to one of the sending tube guides in place of the train of ten pulses. This method of clearing stores is also



TRANSFER UNIT

time of the Dekatron tube. A small overlap has been provided as a precaution against possible failures. If pulses are applied to one guide only, the discharge returns to the same cathode after each pulse.

The method used to pass a digit into the tube is to apply pairs of pulses of the true shown in Fig. 1 to the guide provided to the pulse of the true shown in Fig. 1 to the guide.

The method used to pass a digit into the tube is to apply pairs of pulses of the type shown in Fig.3 to the guide clectrodes. The discharge will then advance by n steps, and if the discharge starts from zero, or output, cathode, it will move to cathode n. To read a digit out of a storage tube, a train of ten pairs

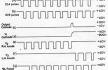
of place may be applied to the tube, causing the discharge pulses may be applied to the tube, causing the discharge pulses may be applied to the tube, causing the discharge produced as the discharge passes the zero or output cathode, and the number of pulses of the train of 10 which occur after the output pulse is an indication of the digit stored on the tube. Further, the number of pulses preceding the output cathode pulse is the complement on 10 of the digit stored in the tube.

The function of the transfer unit is to accept the output coulded paider (nor the Dekatern the form which the three the property of the property of the paider property of the paider preceding the digit and complement, either of which is pared to the econyst table. Fig. 4 is a certail the paider property of the paider of the paider of the paider of the paider property of the paider of the paider of the paider appearing when the digit 3 is transferred. When an addition is to be performed, a refur is energied in a paide to in the paider of the paider of the paider of the paider transferred from the relay circuits to die away, produce transferred from the relay circuits to die away, produce transferred property of the paider of the paider of the transferred property of the paider of the paider of the forest telps round by one complete revolution. When the discharge arrives art to origin cathody we revolved my of the discharge arrives art to origin cathody are very deep or forest telps round by one complete revolution. When the discharge arrives the origin cathody are the paider of the forest telps of the paider of the paider of the paider of the discharge arrives are the paider of the paider of the paider of the paider of the discharge arrives are the paider of the paider of the paider of the discharge arrives are the paider of the paider o

Pulses are applied to the cathode of  $V_s$  synchronously with those shown in Fig. 5(b), but at a different Dc. level, and, prior to  $V_s$  being triggered, they produce negative pulses at the right anode of  $V_s$  (Fig. 5(e)). When  $V_s$  is triggered the pulses appear at the left anode (Fig. 5(fi)), it will be seen that the left anode delivers the direit, and the

used when the machine is first switched on, since the Dekatron tubes may light in any position. When the traisfer operation is carried out purely to clear a store no receiving tubes are connected, but a dummy circuit (the drain address "00") is connected to simulate a receiving address, in order to satisfy the checking arrangements in the address selectine circuits.

To extinguish V, and minimize the chance of spurious triggering taking place, H.T. is removed from its anode from the end of each operation until the start of the next.



As well as the transfer circuits described above, circuits are needed to perform carry over, when required, from a consist of the consist of a cold cathode valve which is riggered when the receiving soler reaches to passe zero. This valve when the receiving soler reaches to passe zero. This valve when the receiving soler reaches to passe zero. This valve pulses, when a second valve performs the carry over and the first is extinguished. A transfer and a carry circuit is the contraction of the contractio

MULTI-DIGIT ADDITION AND SUBTRACTION

circuit, and carry over is made into the least significant

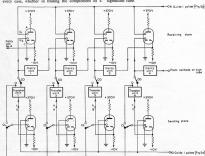
tube. This carry circuit plays no part in handling positive numbers, but is necessary to handle complements. When mumbers, but is necessary to handle complements. When plements on 10 is required for the least significant digit and the complement on 9 for all the other digits. It has arready been explained that in this machine the complement characteristic properties of the properties of the complement of a significant figure and the content of the complement of a significant figure. Carry over from the most significant tube to the least significant corrects this, as can readily be in the complement of a in every case, whether in finding the complement of a in every case, whether in finding the complement of a

all carry units indicates that the carry operation is complete. This is arranged to stop the pulse generator and send a "finish" signal to the relay sequence control circuits, which then proceed with the next step.

which them proceed with the next step.

I can be a proceed with the next step.

I can be a positive number, or a complement (representing a negative number), or a complement (representing a negative number). This could be done by some device (for example a cold cathode trigger valve) having offer the content of the cont



Units Tenths Hundredths
Fig. 6. Arrangement of transfer and curry units for three digit operation.

number or finding the number from its complements.

The partic agencies of the contract provide in particular by the particular provides in Fig. 50 (tappied to receiving subset) after the main in Fig. 50 (tappied to receiving tubbe) after the main contract provides and the particular provides and particular provi

Sign

For positive numbers the sign tube is at zero, and for complements at 9. A double troide is connected to the output cathode of the sign tube, and this operates two relays corresponding to 4 and -. The contacts of these relays give sign information required when multiplying, dividing or printing. It is also arranged that a warring is given tube, perhaps as a result of faulty address selection on MULTIPLEATION SND DIVISION.

Division necessitates the use of a special store (the accumulator) having 15 digits, since when an 8-digit quotient is to be obtained with an 8-digit divisor, there are cases where difficulties would arise if a 15-digit divisor were not provided. An accumulator of more than eight digits also reduces errors in rounding off the product of multiplication if products are repeatedly summed in the accumulator. A further advantage of a long accumulator.

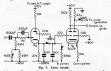
is that the remainder after division is available for use if

required.

In both multiplication and division ose storage addres is used as the sending store, the accumilation of the storage address is used in a special way, to be described, and is referred to as a register. The contents of these stores are stage of each of these stores and each of the store and the stores are stage of each of these stores and tensor that the store and the stores and tensor that the addresses have been selected relays connected to these circuits of these circuits of the stores and trained that the store and trained to the store and trained that the store and the stores are stored to the store and th

TABLE I		
	MULTIPLICATION	DIVISION
SENDING STORE ACCUMULATOR	MULTIPLICAND PRODUCT MULTIPLIER	DIVISOR DIVIDEND OLOTHENT

Where the multiplier is positive, the multiplicand is added into the accumulator V, times where V, is the number to the accumulator V, times where V, is the number to the control of the



it is actually moved on nine steps without carry over) for each single addition. The pulse generator is arranged to give the finish signal calling for the next operation when the appropriate digit of the register has reached zero and earry over is complete.

When the multiplier is a complement, a similar procedure is carried out with a few modifications. Multiple subtractions are required, and the register tube is moved forward one step for each subtraction. Since the complement of 0 is 9, it follows that transfers should stop when the register tube reaches 9. As the pulse generator allows transfers to continue until the tube reaches 0, one too many will be performed by the time the finish signal is given, and it is necessary to make a single addition to correct for this. The sequence then is multiple subtraction, single addition, operation of shift. These sequences are controlled by relays in the sequence controlling and routing section of the machine. It is, of course, only necessary to give the order "multiply"; the precise sequence required is selected before the operation starts by a relay circuit dependent on the sign of the contents of the selected addresses

Division is performed by a sequence of the same form,

i.e., multiple transfer, single transfer, operation of the shift; the quotient is built up digit by digit in the register by moving the discharge one step forward or backward in the register for each transfer. For a divisor and dividend of the same sign, the divisor is subtracted from the dividend and the register moved forward until the sign of the dividend changes, when the pulse generator gives the finish signal. At this point one subtraction too many has been performed, and this is corrected by making one addition and moving the register back one step. When the divisor and dividend are of opposite sign, the multiple transfers are additions, with the register being moved back step by step and the single transfer is a subtraction. In this case it is not necessary to move on the register during the single transfer, since the fact that it has started from 0 rather than 9, the correct starting point for a complement, means that the necessary correction has already been made. ROUNDING OFF

After multiplication or division, some form of rounding off is normally needed to avoid a systematic error due to the omission of digits less significant than the least significant held in the stores. In multiplication, rounding off is most conveniently done when the product is transferred from the 15-digit accumulator to an 8-digit store. A method commonly employed is to add five to the most significant of the digits which are omitted in this process; carry over then adds one to the least significant figure retained if the omitted figure was 5 or over. The maximum error in the last digit is +0.5, and the average error is zero. This process would have proved inconvenient in the present machine, partly because it would involve the generation of a special train of 5 pulses, but also because in the case of division, the quotient is built up in an 8-digit store, and without much extra complication a 9th digit is not available to be rounded off. Accordingly the cruder method of adding I or 0 at random to the 8th figure of the product or quotient has been adopted. The maximum error in this digit is now ±1, and the average error is zero. The probable error after rounding off in this way is about twice that to be expected from the more conventional process. This error is not considered serious. A more important point is that the random nature of the result makes difficult the precise checking of the machine by repeating a solution in which rounding off had taken place.

To minimize systematic errors the selection of 10 or 1 for rounding off must be as nearly random as possible. The method adopted is relatively simple. A muster oscillator which are subsequently routed to perform the transfer and carry operations. Pulses from this oscillator are led controlled to the control of the control of the control tricks which conduct alternately. When rounding off into any address the state of the scale-of-two circuit at 0 to beginning off the transfer period determines whether the

The relation between the state of the scale of two at the beginning of one round off operation and that at the beginning of the previous round off that depends upon whether the number of pulses occurring in the interval whether the number of pulses occurring in the interval miliseconds apart whereas the interval between successive round off operations is several (usually very many) seconds. The exact length of this interval is affected by many uncertainties, such as the operating interval of the tape readers of the tape readers of the previous control of the product of the previous control of the pr

Experiments so far carried out show no significant correlation between the zeros and ones produced by repeated round off operations.

The Translator Unit

The function of the translator unit is to convert the

information obtained via a relay circuit from a punched paper tape into a form suitable for operating the transfer units in the computor and vice versa. The translator is brought into circuit in the same way as any storage address and no alteration is needed to the transfer unit to accom-

The basis of the unit is a special Dekatron in which the cathodes are all brought out separately. For this purpose a 10-cathode tube would be suitable, but the tube available has 11 cathodes (GS11A) and this has been used.\*

The method of operation is shown in Fig. 8. If, for

the property of the property of the property of the period of the period

on which the discharge is left rises in potential from -16V to something positive with respect to earth, causes the grid potential of two of the triodes VII to VI5 to rise, and so closes two of the relays Pa(1, PB/1 - - PE/1, The contacts of these relays operate the teleprinter or tape perforator. The metal rectifiers in the cathode circuits of the CS1/A give an output in the appropriate code (i.e., success) for this purpose.

After translation-in or -out a separate circuit returns the discharge in the GS11A to cathode No. 0 before the translator is again required for use.

### Sequence Control

This part of the computor controls the tape readers and output organs, selects the Dekatron stores required for each arithmetical order and interprets these orders as a sequence of the basic functions of addition and subtraction. It is designed to operate as an asynchronous system in which each switching operation in the sequence is

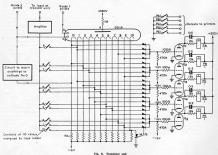


Fig. 8. Translater unit

of the train of 10 (in the example, 3) are passed to the checked before the next is allowed to begin. This simplifies

of the train of 10 (in the example, 3) are passed to the receiving store by the transfer unit. When translating in negative numbers the transfer unit is arranged to complement them. Since the translator deals with only one digit at a time the remaining transfer circuits are arranged to transfer zeros for nines in the case of a complement). The translator is connected via the transfer circuits to each dieit of the store in turn by the relay shift circuit.

When translating out (i.e., taking information from the competior) the GS11A tube again starts from cathode No. 0, but in this case it steps round by the number of pulses delivered to it by the transfer unit. The cathode wayeneds GS11A, the cathode wayeneds GS11A and 12 cathode, which wayeneds GS11A.

maintenance by eliminating critical relay timing and by stopping the computor immediately a circuit fault is detected.

The sequence of an arithmetical operation can be regarded as consisting of three parts:—

(1) Reading in the current order.

Setting up the necessary connexions.
 The arithmetical operation itself.

Non-arithmetical orders are dealt with in a similar manner and the operation of the computor continues until a special order is reached which leads to the lighting of the "Finish" lamp and prevents any further operations. READING IN THE CURRENT ORDER

Each order is transferred, digit by digit, from the perforated tape or Dekatron store into the relay order storage circuit and is then interpreted as a sequence of additions and subtractions. The address from which the order is obtained has been derived from an earlier order calling for a change of order source. This address is stored on a group of relays which mark out the two digits of the order source address. The "tens" digit energizes a relay connecting a group of ten addresses as the sending store and the "units" digit selects the individual store in this group. The reading-in sequence begins when it has been proved that the correct tens and units relays have operated. If the order source is a tape reader the sequence will consist of six cycles. The space or block-marking character is read during the first cycle and the order digits during the remaining five cycles. At the beginning of each cycle tha appropriate relay storage group is connected to the five-

wire output of the tape readers. The contacts of each tape reader correspond with the pattern of holes on the tape, and they operate the equivalent combination of five relays in the storage group to which they have been connected. The relay storage group into which the space or block-marking character should be routed will only allow the next cycle to commence if it receives a particular combination of one hole out of five (the space combination) or any combination of three holes (the block-marking characters). The five storage groups into which the order digits are fed will only allow the next cycle to commence if each receives one of the ten combinations of two holes. If the tape reader moves the tape forward more than one row before reading the next character, or fails to move the tape, a digit storage group will be offered the space code or vice versa. Furthermore, any erratic operation of the reading contacts is unlikely to distort a two-hole code into another twohole code, so that any faulty operation of the tape reader is almost certain to be detected.

SETTING UP THE CONNEXIONS

The source from which the current order has been read and the relay storage group are now disconnected, leaving the order stored on the relays. If this order requires an arithmetical operation the sending and receiving addresses are selected and connected to the input and output of the transfer system in the following way. The Dekatron stores are built in units each containing ten stores of eight digits and sign. These units are a convenient size and suit the decimal address code. Each group of ten stores has switching relays which connect it either to the input or the output of the transfer systems. These are the sending and receiving "tens" relays. When a "tens" relay is energized the Guide 2 electrodes of all the sign storage Dekatrons in the group are connected to the sign input or output lead of the transfer system, all the Guide 2 electrodes of the 1st digits to the 1st digit lead, and so on. The individual store in the group is selected by one of ten switching relays (the "units" relays) which connects Guide I pulses to all the Dekatrons of one address. Thus although all the 90 Dekatrons in the group of 10 addresses receive Guide 2 pulses, only those of the selected address receive pulses on both guides and the position of the discharge can change only in this one address.

The two sets of "tens" and "units" relays which

bring the sending and receiving addresses into action are selected and checked by four groups of address storage relays. A fifth group stores the operation to be performed and, provided that the addresses selected are compatible with the operation demanded, sets up the arithmetic control relay circuits. As a further guard that the sending and receiving addresses are connected correctly the arithmetic sequence cannot start until the signs of both addresses and of the accumulator have been identified.

THE ARITHMETIC OPERATION When the selection of function and addresses is complete the arithmetical sequence begins. It occupies a variable number of cycles, only one for addition or subtraction, nine for multiplication, ten for division, 11 for reading in a number from the tape and up to 19 for printing a number in a maximum length print layout. For addition or subtraction the single cycle has two steps; setting the shift circuit to the straight-through condition and carrying out a single transfer. For division of a positive dividend by a positive divisor each cycle has three steps: setting the shift to successive positions, carrying out a multiple subtraction until the sign of the accumulator changes, and carrying out a single addition. The comple-

tion of each step is checked before the next can begin. When the arithmetic sequence has produced the full number of cycles required for the particular operation the address relays are released, the stored order is cleared down and the next order reading sequence begins.

PRINTING OUT If the current order selects the printer as the receiving address the translator is connected to the output of the transfer system and the printer is connected to the fivewire output of the translator. A previous order will have specified the print layout which is to be adopted. Each cycle of the arithmetical sequence then consists of three The first sets the shift to route each digit in turn to the input of the translator (which accepts only one digit at a time). At the same time the translator is recycled to zero. During the first step there may alternatively be connected to the printer the contacts of one of a number of relays which mark out non-digital characters such as space. line feed and carriage return. The second step is a single addition or subtraction. If a non-digital character has been selected by the first step this transfer is not required. The teleprinter with page printing attachment has been modified so that five-wire signals operate magnets controlling the five combs of the combination head. The third step of each cycle sets the combination head of the teleprinter by operating these magnets from the output of the translator or a character relay, checks the position of the combs, and prints a character.

Conclusion

It is believed that the computor described above will meet a real need in the computing field. It is particularly suited for cases in which the same calculation must be performed many times on a variety of numerical data. Its use is expected to be an economical proposition in many cases when the cost of a large, fast computor would not

be justified Acknowledgment

The authors have received much beloful advice from Dr. J. Howlett, and during the early stages of the work they benefited by a discussion with Professor D. R. Hartree. The work was carried out as part of the programme of the Electronics Division at A.E.R.E. and is published by permission of the Director, A.E.R.E.

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