ELC-Europe 2009 - Grenoble

Philippe GERUM - SourceTrek

# Introduction

#### The context

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  - PREEMPT\_RT close to mainline

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- Linux about to be natively real-time
  - PREEMPT\_RT close to mainline
- Legacy applications knocking on Linux's door
  - Traditional, embedded RTOS
  - Non-POSIX core API
  - Flat / physically addressed memory
  - Typically: VxWorks, pSOS, VRTX etc.

#### The issue

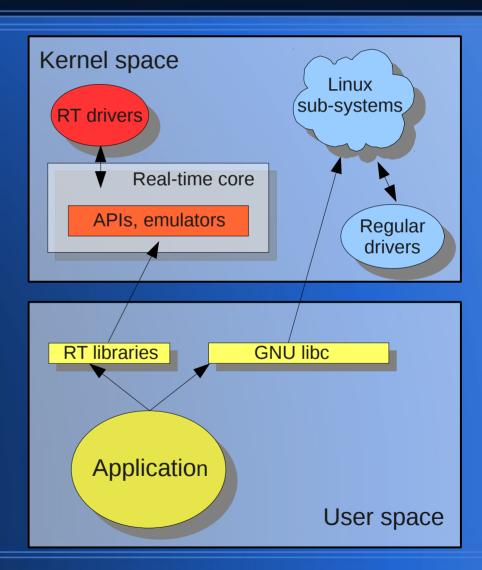
- Porting them to Linux currently means
  - Rebasing on Linux, changing design or,
  - Keeping design, keeping proprietary RTOS
- How to go the Linux way?
  - Keeping design, using Linux technologies

#### Possible solution

- Combine existing Linux technologies
  - Native real-time support
  - Linux-native virtualization
  - RTOS emulation

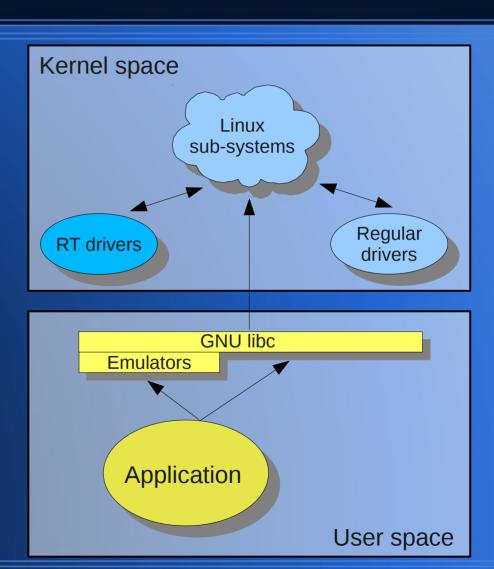
#### Common porting strategies

- Port to dual kernel
  - Over POSIX API
  - Over API emulator
  - Over ad hoc API

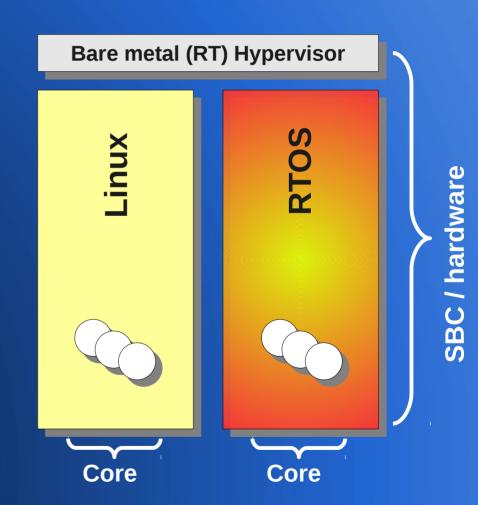


#### Common porting strategies

- Go Linux native
  - Over POSIX API
  - Over API emulator

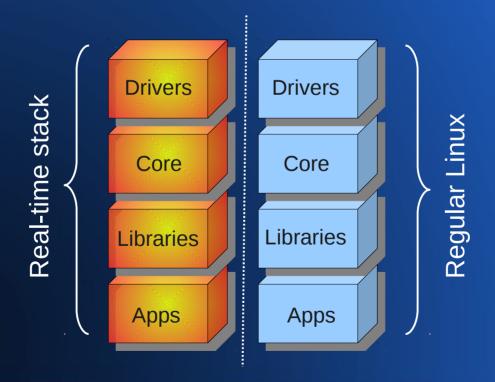


- Introduce virtualization
  - Original RTOS guest
  - Vendor-specific
- Bare metal hypervisor
  - Leverage multi-core

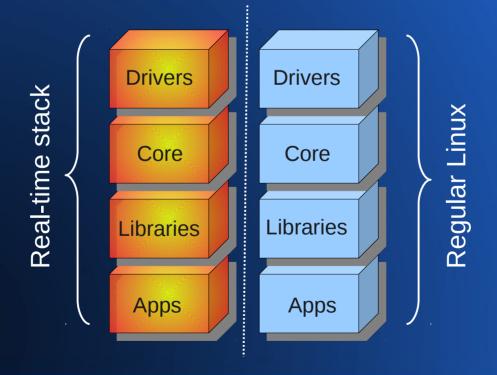


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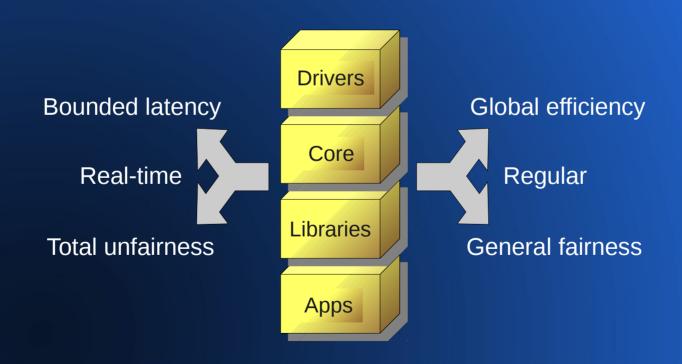


- Dual kernel Linux architecture is complex
  - Pressure on application design

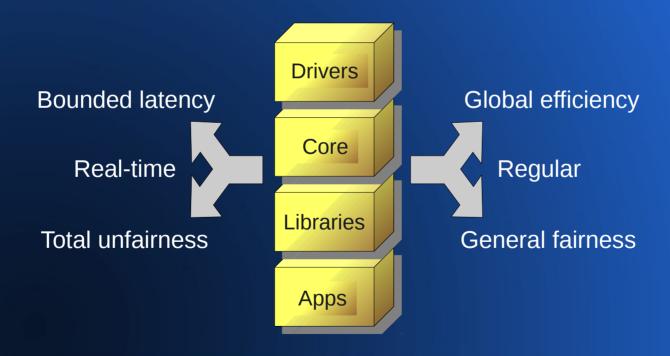


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- Native real-time Linux is complex
  - Pressure on system configuration



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    - Hypervisor technology
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    BUT,
  - Do not help the Linux real-time effort

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  - Over-emulation of missing calls
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- Driver model
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  - Linux kernel API is more complex
- Protocol stacks
  - Keep "as is" or offload to Linux?

#### Legacy issues

- Software architecture
  - BSP code exposed
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  - Non-public API sometimes used

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- Software architecture
  - BSP code exposed
  - Application and driver code entangled
  - Non-public API sometimes used
- Programming model
  - Flat / physically addressed memory assumed
  - Supervisor mode assumed
  - CPU architecture assumed

# About RTOS emulators

#### RTOS API emulation?

- A way to mimic the RTOS interfaces
  - Evades the BSP issue
  - Source-level approach
- Has real-time requirements
  - Must run over a deterministic core
  - Must exhibit real-time properties itself

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**YES** 

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Do you trust your vendor documentation? YES

Should your code <u>rely on undocumented</u> features? **NO** 

Should your code expect undocumented behavior? NO

Therefore, you don't need the original API implementation to emulate it properly.

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  - Traditional RTOS share basic semantics
    - Optimized building blocks can be made
    - Efficient "window-dressing" follows
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  - Traditional RTOS share basic semantics
    - Optimized building blocks can be made
    - Efficient "window-dressing" follows
    - Leveraging single address space helps
  - Naive emulation over POSIX not enough
    - POSIX semantics do not map 1:1
    - POSIX-based <u>building blocks</u> may work better

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  - noarch/generic core services
- Require Application / Driver split
  - BSP code not accessible from user-space
  - I/O resources live in kernel space
- Restricted by Linux protections
  - No supervisor actions from user-space

# Our assets

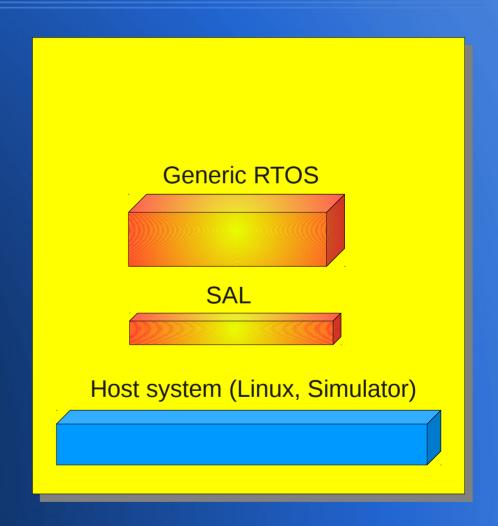
### PREEMPT-RT

- Fully native real-time support
  - Enables real-time virtualization
- Promise of embedded multi-core scalability
  - Sophisticated locking model
  - Sophisticated scheduling

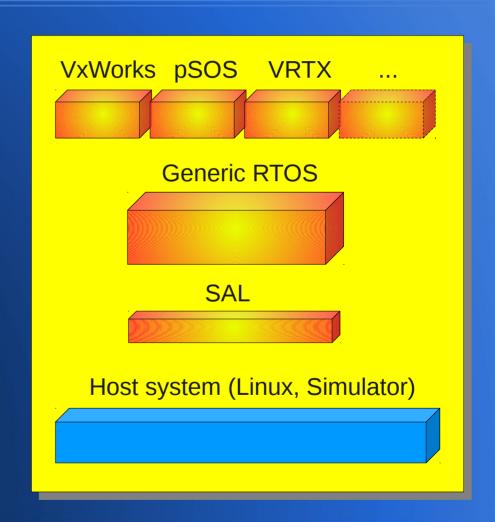
### **KVM**

- Complete sandboxing
- Compatible memory spaces
- Device virtualization through host
  - virtio
- Device emulation through VM
  - Qemu-based modelling

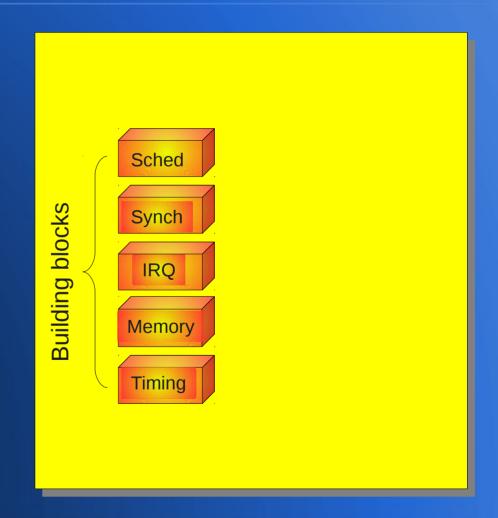
- Generic RTOS core
- Host abstraction
  - Dual kernel
  - Simulator
  - (Single image \*)



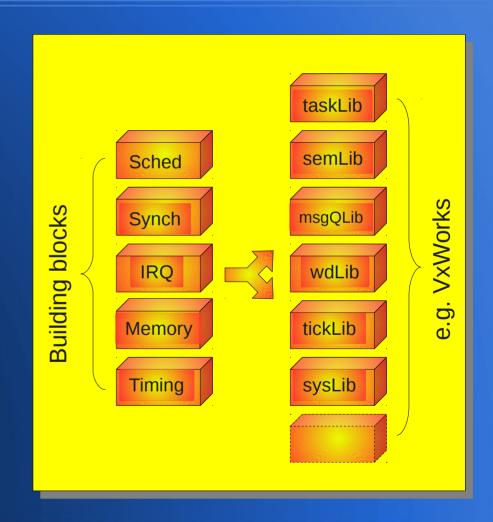
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- RTOS personalities



- RTOS building blocks
  - Thread scheduling
  - Synchronization
  - Interrupt handling
  - Memory allocation
  - Timing services

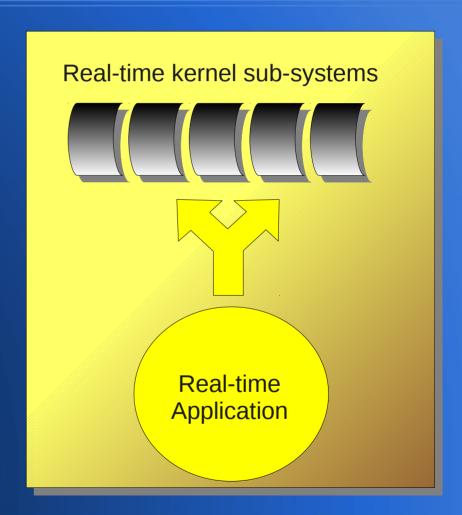


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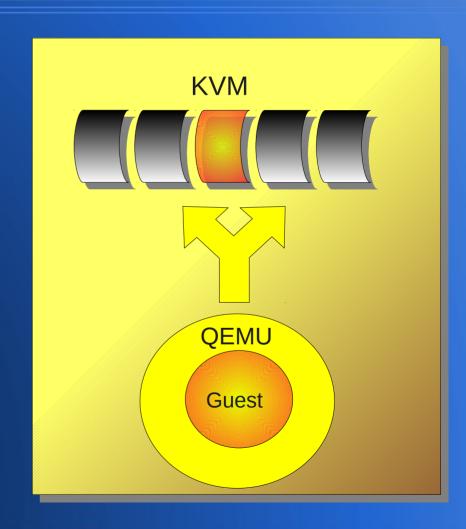
# What about combining?

- Real-time host kernel
  - PREEMPT-RT



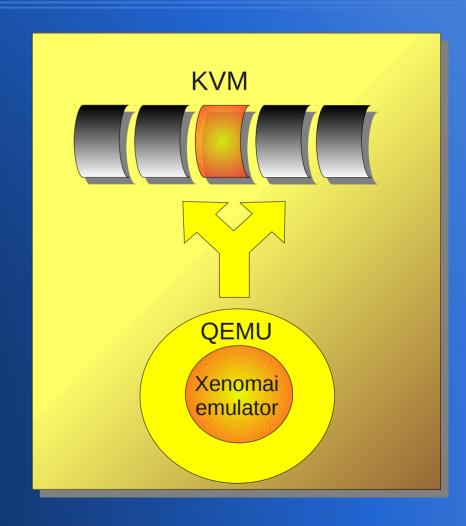
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  - KVM
  - QEMU



## What about combining?

- Real-time host kernel
  - PREEMPT-RT
- Virtualization core
  - KVM
  - QEMU
- RTOS emulation
  - Xenomai



### Virtualization + RTOS emulation

### **Improvements**

- Native real-time
- Original programming model
- Better emulation coverage
- Sandboxing
- Legacy device emulation

### Virtualization + RTOS emulation

### Restrictions

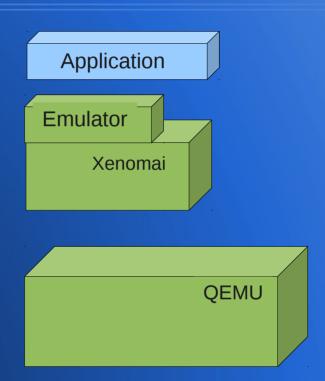
- No ABI compatibility
- Still not 100% source compatible
- Reworking the device driver layer still required

# Virtualize & Emulate

## Improved emulation engine

### **Emulation core**

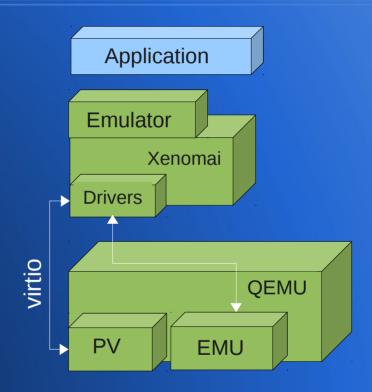
- Xenomai guest
  - Freestanding mode
  - RTOS personality
- QEMU
  - Virtual machine



## Improved emulation engine

### Handling I/O

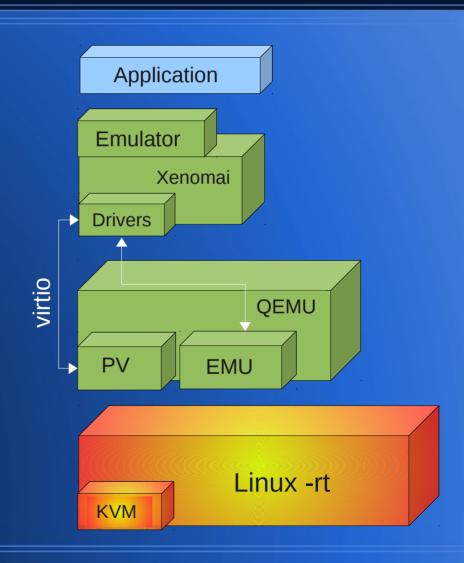
- Paravirtualized
  - Common hw
  - High bandwidth
- Emulated
  - Precise emulation
  - Low bandwidth



# Improved emulation engine

### **Native real-time VMM**

- PREEMPT-RT host
  - KVM-enabled



### **TODO list**

- Real-time aware KVM
  - Guest scheduling
- Real-time aware QEMU
  - I/O emulation
- Guest mode Xenomai core
- Extended emulation coverage

# More applications

### Could also be used for...

- Application-specific virtual RTOS
  - Virtual RT appliance (sort of)

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- Transition path for in-house RTOS
  - Consolidate & extend via virtualization

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- Application-specific virtual RTOS
  - Virtual RT appliance (sort of)
- Transition path for in-house RTOS
  - Consolidate & extend via virtualization
- Simulation of complex architectures
  - e.g. modeling Arinc653 systems

# Conclusion

# Legacy RT application to Linux

### **Today**

- Rebase on Linux, change design
- Keep design, keep proprietary RTOS

# Legacy RT application to Linux

### **Today**

- Rebase on Linux, change design
- Keep design, keep proprietary RTOS

### **Tomorrow**

- Combine existing technologies
  - Rely on real-time capable virtualization
  - Couple with accurate RTOS emulation

### The End

Thank you for attending