Embedded Alley





Memory A Most Precious Resource

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Introduction

Why Memory?

Capacity isn't time-based

If you need more, someone has to free

CPU too slow, it just takes longer

What to do?

Protect my stash

Ask Linux OS for assistance





How Do We Do It?

- Cgroups Overview
- Memory Cgroup
- Memory Usage Notification
- Memory Cgroup Example
- Out of Memory (OOM) Killer
- Future Directions



Cgroups Overview

- Referenced by several terms
 - Containers (or Container Group)
 - Control, Controller Group
- We'll just say "cgroups"
 - What are cgroups?
 - How do they work?
 - Why do I care?



What Are Cgroups?

- A mechanism for partitioning sets of tasks
- Managed in a mounted virtual file system
- Can create a group hierarchy
 - Multiple groups within a tree
 - Root of tree is management point for sub-trees
- > No system cost if not used, insignificant when used







How Do Cgroups Work?



- A resource controller is the policy
- Tracks the reference counted objects
 - cpu usage
 - cpu sets
 - memory pages
- Works with a kernel subsystem for resource management



Why Do I Care?

- Powerful system resource management concept
- Resource consumers become part of the management
- OS doesn't have to guess (sometimes poorly)
 - Thread priorities are just hints
 - VM tuning knobs can be a research career
- Tuning moves to the application space
 - Ill-behaved tasks in their own container
 - Easily accommodate feature enhancement
 - Sensible system tuning perspective





- Mounted virtual file system (i.e. /cgroups/<restype>)
 echo pid > /cgroups/<restype>/<userclass>/tasks
- Easily and dynamically change resource controls
 echo pid > /cgroups/<restype>/<new_class>/tasks
- Threads can determine cgroup information
 - /proc/<pid>/cgroup



Memory Cgroup

- Memory Resource Controller
 - Don't confuse with hardware memory controller
 - Track RSS and page cache pages, swap cache option
 - Reclaims through cgroup LRU
- Isolates memory behavior to a group of tasks
 - Prevent "memory hungry" tasks from consuming entire system memory resource
 - Control memory consumption for virtualization
 - Provide a protected container for critical embedded tasks





- Previous, stand-alone /dev/mem_notify
- New approach builds upon the memory cgroup resource tracking
- Kernel configuration option (CGROUP_MEM_NOTIFY)
- Select a cgroup usage limit notification percentage
 - Percentage rather than absolute value
 - Prevents need to update if cgroup is resized
- Task operates normally until notification
 - Can block-wait until limit
 - Can poll as part of normal processing



Memory Notification Cgroup Example

- Create a virtual file system
- Set resource limits
- Memory allocation code fragment
- Notification thread code fragment
- Example program message output



Create Virtual File System

- mkdir -p /cgroups/memcg
- mount -t cgroup none /cgroups/memcg -o memory
- > mkdir /cgroups/memcg/0
- echo \$\$ > /cgroups/memcg/0/tasks



Set Resource Limits

- # ls /cgroups/memcg/0
 memory.failcnt
 memory.force_empty
 memory.limit_in_bytes
 memory.max_usage_in_bytes
 memory.notify_limit_lowait
 memory.notify_limit_percent
- memory.notify_limit_usage memory.stat memory.usage_in_bytes notify_on_release tasks

- Set the memory usage limit
 - > echo 10M > /cgroups/memcg/0/memory.limit_in_bytes
- Set the notification limit to 80%
 - > echo 80 > /cgroups/memcg/0/memory.notify_limit_percent





Memory Resource Notification Example

- Multi-threaded application
 - Main thread allocates 10 segments
 - Main thread frees segments if they still exist
 - Continues in a loop
- Notification thread
 - Blocks on memory.notify_limit_lowait
 - Frees allocated segments until usage < limit</p>



Memory Allocation Code Fragment

```
k = 10;
while (k-- > 0) {
     for(i = 0; i<NSEGS; i++) {
          if ((mp = malloc(SEGSIZE)) == NULL) {
               perror("malloc");
               exit(2);
          memptr[i] = mp;
          for (j = 0; j < SEGSIZE; j++)
               *mp++ = j;
          printf("Alloc seg %d\n", i);
          sleep(5);
     for(i = 0; i<NSEGS; i++) {
          if ((mp = memptr[i]) != NULL)
               free(mp);
          printf("Free seg %d\n", i);
          sleep(5);
```



Notification Thread Code Fragment

```
for (;;) {
     /* Open/read /cgroups/memcg/0/memory.notify limit lowait
     * Blocks while usage is below limit
     */
     percent = get_memcg_val("lowait");
     limit = get memcg val("percent");
     printf("Notify wakeup percent %d, limit %d\n", percent, limit);
     i = 0:
    do {
          if ((mp = memptr[i]) != NULL) {
               memptr[i] = NULL;
               free(mp);
               printf("Notify free seg %d\n", i);
          i++:
          usage = get_memcg_val("usage");
     } while ((usage > limit) && (i < NSEGS));
```



Example Program Message Output

Alloc seg 0 Alloc seq 1 Alloc seq 2 Alloc seg 3 Alloc seq 4 Alloc seq 5 Notify wakeup percent 80, limit 80 Notify free seg 0 Alloc seq 6 Notify wakeup percent 80, limit 80 Notify free seg 1 Alloc seg 7 Notify wakeup percent 80, limit 80 Notify free seg 2 Alloc seg 8 Notify wakeup percent 80, limit 80 Notify free seq 3 Alloc seg 9 Free seg 0 Free seg 1 Free seq 2 and continues



- Moving a task doesn't migrate old allocations
 - Reclaims will deplete old cgroup allocation
 - New allocation charge to new cgroup
- Notification doesn't carry information
 - Normally, wake up is due to reaching notify limit
 - Wake up on thread migrate to new cgroup
 - Cgroup is forced empty (memory.force_empty)
 - Task must interrogate cgroup state to determine action





- Linux chosen method of managing memory overload
 - Often nothing to kill in an embedded system
 - Difficult to make the right choice
- Operation is based upon kernel tuning
 - Memory overcommit
 - OOM adjustment knobs (per process)
 - Policy choices always under discussion
- Memory cgroup is subject to OOM
 - Overload will trigger OOM within cgroup
 - Can leverage OOM cgroup (http://lwn.net/Articles/315949/)



Future Direction

- Improve notification API
- Additional cgroup subsystems
- Increased granularity
 - requires lower overhead
 - asynchronous notification
 - information arguments passed with notification
- Active resource management programming model
 - Application states "...this is what I plan to need ..."
 - At end of block the need is revoked
- Find some assist for legacy applications



Summary

- All system resources are precious and must be managed
- Cgroups provide the mechanism for task partitioning
- A subsystem resource controller provides the policy
- Enables a powerful application-centric resource management
- Memory notification patch is in the e-mail queue to lkml