

Practical Filesystem Security

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Hello

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- > Focus on Linux kernel, low-level components, virtualization, security

Overview of this Talk

- > Practical overview of filesystem security on embedded Linux systems
- > Hopefully some guidance for your next project
- > By no means a complete guide how to implement a whole security concept
- More a collection of pointers

Motivation for Filesystem Security tongue-in-cheek

- > Care about customer data on the device
- Care about data integrity
- › Keep your magic sauce secret
- Have creative licensing
- > Pass some certification test

Know your threat model

- > Has attacker hardware access?
 - > To a running device?
 - > Able to dump main memory?
- Access to a shell?
 - root?
 - kernel level?
- Who is the attacker?
 - Nosy neighbor?
 - Competitor?
 - Secret agency?

Filesystem encryption

- > Encrypted...
 - Disk?
 - > Filesystem structures?
 - Files?
 - Directories?
 - > File names?
 - > Out of band data (xattr, ...)?

Filesystem encryption: eCryptfs

- › Kernel mode stacked filesystem (no FUSE)
- > Encrypts file content and file names on top of another filesystem
- Per directory basis
- No authenticated encryption

Filesystem encryption: Possible eCryptfs issues

- > Performance overhead from stacking
- > File name limit

```
$ stat -f -c "maxlen: %1" /some/ecryptfs
maxlen: 143
$ stat -f -c "maxlen: %1" /other/fs
maxlen: 255
```

> Who of you checks file length limit before creating a file?

Filesystem encryption: Possible eCryptfs issues (cont'd)

- > On Linux a file name must not contain a nul byte or a slash
- > Encrypting a string can give you any result, including nul bytes or slashes
- > eCryptfs has to encode cipher text: Increases length

Filesystem encryption: Using eCryptfs on Yocto

- Add ecryptfs-utils to your rootfs
- > Enable CONFIG_ECRYPT_FS in kernel config
- > mount ecryptfs before you need it
 - initramfs
 - > PAM
 - application level

Filesystem encryption: dm-crypt

- > Block level encryption, uses device mapper
- Works with any block based filesystem
- Used for FDE (Full Disk Encryption)
- > Rich cipher suite
- No authenticated encryption

Filesystem encryption: Using dm-crypt in Yocto

- > Add cryptsetup to your rootfs
- > Enable CONFIG_DM_CRYPT in kernel config
- > Setup dm-crypt before you mount the filesystem
 - > Happens usually in initramfs

Filesystem encryption: fscrypt

- > File encryption at filesystem level (no stacking)
- > Currently supported by ext4, f2fs and ubifs
- > File content and file names are encrypted
- No meta data nor out of band data (xattr)!
- > Per directory basis (per directory encryption policy)
- Per inode AES key
- No authenticated encryption

Filesystem encryption: fscrypt (cont'd)

- > Primary use case: per user and directory encryption
- > Can be abused to encrypt whole filesystem
- Master key provided via keyctl
- > Key has to reside in an accessible keyring (e.g. session keyring)
- > Has no problem with long file names.
 - > Quiz question: Why doesn't it suffer from the same problem as eCryptfs?

Filesystem encryption: Possible fscrypt issues

- > pam_keyinit is your enemy
- > File content in page cache, if user A has a key and reads a file, user B can read it too if access control allows it!
 - > Consider mount namespaces or strict DAC/ACL
- > Without the key nobody can read cipher text
 - No backup possible without key!

Filesystem encryption: Using fscrypt in Yocto

- Add fscryptctl to your rootfs
- > Enable CONFIG_FS_ENCRYPTION in your kernel config
- > After mounting fileystem make sure either all or selected users have a key

Filesystem encryption: More considerations

- > Full disk encryption is the last resort option
- > Think of fine grained encryption, eCryptfs or fscrypt help here
 - > Do you really need an encrypted /usr and /lib?
 - If possible, combine dm-crypt and eCryptfs/fscrypt

Filesystem encryption: What about data integrity?

- > Changed ciphertext usually remains unnoticed
- › Just decrypts to garbage
- > Attackers can still do evil things
- > Think of block swapping or swapping whole (encrypted) files
- > e.g. if location of true and login are known their content can get swapped
 - No plaintext needed
- > Pre-generated filesystem images help attackers

Filesystem integrity: dm-verity

- > Read-only device mapper target
- > Useful for read-only block based filesystem such as squashfs or erofs
- > Fast, uses a hash tree
- Use cryptsetup/veritysetup on target
- > CONFIG_DM_VERITY in kernel config

Filesystem integrity: dm-integrity

- > Read-write device mapper target
- > Basically adds an auth tag to every block
- Can be combined with dm-crypt
- Use cryptsetup/veritysetup on target
- > CONFIG_DM_INTEGRITY in kernel config
- › Non-negligible overhead

Filesystem integrity: fs-verity

- > Integrity for selected files
- Read-only!
- > Supported on ext4, f2fs and btrfs
- Use fsverity-utils on target
- > Enable CONFIG_FS_VERITY in kernel config

Filesystem integrity: authenticated ubifs

- > Full authentication support and read-write
- > Works because ubifs is strictly copy-on-write
- Can be combined with fscrypt
- > Be aware: Featre is rather new

Wait, what about generating images?

- > Most mechanisms don't have tooling to generate encrypted/authed images
- > We don't recommend it
- > Installer approach:
 - > rootfs as tarball
 - Generate an installer (IOW a livecd)
 - > The installer will setup everything, plus locking down the device

Filesystem encryption: Using fscrypt with mkfs.ubifs

- mkfs.ubifs can generate a pre-encrypted ubifs filesystem, whole filesystem same policy
- > mkfs.ubifs -r rootfs/ -m 2048 -e 126976 -c 1024 -o
 ubifs_crypt.img -b ddeeaaddbbeeeeff -K ubifs_masterkey.bin
- > ddeeaaddbbeeeeff is the key descriptor, see fscryptctl

Filesystem integrity: Using ubifs authentication with mkfs.ubifs

- Just like for fscrypt
- > We use the signing key from the kernel build
- > mkfs.ubifs --hash-algo=sha256 --auth-cert=signing_key.x509 -r
 rootfs -e 126976 -o ubifs_auth.img -c 1024 -m 2048
 --auth-key=signing_key.pem

The "key" to success

- > No human interaction wanted (e.g. mount must not ask for passwords)
- > Key material must be stored in device itself to unlock device
- > Attacker must not extract key
- Major challenge
- > No way without support from hardware

The "key" to success: Naive approach

- > Derive key from hardware properties
- > CPU ID, MAC from network card, etc...
- > Security by obscurity, IMHO
- More often used than you'd assume

The "key" to success: External Secure Element (TPM, etc.)

- > Can store key material in a secure way
- > Problem: Doing all crypto on the secure element is slow
- > To utilize CPU, key needs get transferred into main memory
- > Attacker can read the key while it is transferred
- Common attack: Bitlocker TPM sniffing

The "key" to success: Internal Secure Element (i.MX CAAM, etc.)

- Some SoC have a built in secure element
- › e.g. i.MX CAAM or DCP
- In short: SoC can do AES with a fused key
- > Typical use case: Store encrypted FDE key on distrusted location
- > Problem: Fails if attacker can execute code, you need verified boot
 - > Applies to the external secure element case too

The "key" to success: Key not in main memory

- Common requirement: KEY MUST NO RESIDE IN RAM!!!11elf
- > Technically possible if you have a secure element
- Keep in mind:
 - Some mechanisms need the key in plaintext and do manual key derivation, e.g. fscrypt
 - > Linux's page cache is not your friend
 - > Consider RAM encryption too
 - Know your threat model!

A few words on performance

- Crypto on SoC can be slow
- > Crypto accelerators are not always faster
 - > Filesystem encryption/auth is not their use-case
- Consider using AES-128 instead of AES-256
- > When using dm-crypt, consider no-read-workqueue and no-write-workqueue
- › Do your own benchmarks!

Summary

- > Know your threat model
- > There is no one-fits-all solution
- > Know your threat model
- > Full disk encryption is the last resort
- > Know your threat model
- > Storing the key material is the hard part
- > Know your threat model

Further reading

- https://www.kernel.org/doc/html/latest/admin-guide/device-mapper/dm-crypt.html
- https://www.kernel.org/doc/html/latest/admin-guide/device-mapper/verity.html
- > https://www.kernel.org/doc/html/latest/admin-guide/device-mapper/dm-integrity.html
- https://www.kernel.org/doc/html/latest/filesystems/fscrypt.html
- https://www.kernel.org/doc/html/latest/filesystems/fsverity.html
- https://www.kernel.org/doc/html/latest/filesystems/ubifs-authentication.html
- https://www.spinics.net/lists/linux-mtd/msg08477.html
- https://www.jakoblell.com/blog/2013/12/22/practical-malleability-attack-against-cbc-encrypted-luks-partitions/
- https://pulsesecurity.co.nz/articles/TPM-sniffing
- https://blog.cloudflare.com/speeding-up-linux-disk-encryption/

FIN



Thank you!

Questions, Comments?

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