Debugging Real-Time issues in Linux

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Life of an RT engineer
Real time - terminology

For purposes of this task:

Period: Time Interval between which RT task will be released at fixed rate

Deadline: Time from when an event occurs to when RT tasks finishes response

For periodic Real time tasks like audio, Period = Deadline
Fig. (1a) - Example of working periodic real-time tasks with deadlines equal to periods. Not deadline misses.

Fig. (1b) - Example of problematic periodic real-time task with deadline misses due to too long execution time.

Fig. (1c) - Example of problematic periodic real-time task with deadline misses due to too jitter (scheduling latency).

Fig. (1d) - Example of problematic periodic real-time task with deadline misses due to too jitter combined with too long execution time.

Image Source: Eduardo Valentin <edubezval@gmail.com>
Life of a Wake up (Interrupt 2, Task 2 responds)

Delays in Task 2 response to an event

CPU interrupts masked

Interrupt 1 ISR

Interrupt 2 ISR

Task 1 (higher priority process)

Task 2 gets a chance to run

Hardware delays

IRQS off delays

ISR execution delays

Scheduler delays

Task response delays

Time

Interrupt 2 Signal

CPU receives signal from IRQ controller

Task 2 is woken (placed on CPU run queue)

Task 2 is given CPU

Response complete

Total Response Delay
Life of a Wake up with Preempt Off delays

Delays in Task 1 response to an event (Task 1 is higher priority than task 2)

- Task 2 is running
  - Task 2 - Preemption turned off
    - Interrupt 1 ISR
      - Hardware delays
      - Interrupt Servicing delays
      - Preemption Off Delays
      - Task execution delays
          - Task 1 gets a chance to run
            - Task 2 is running
                - Task 2 - Preemption turned off
                    - Interrupt 1 Signal
                      - CPU receives signal from IRQ controller
                        - Task 1 is woken (placed on CPU run queue)
                          - Preemption turned on, Task 1 is given CPU
                            - Response complete

Total Response Delay
Life of a Wake up on an IDLE CPU

Delays in Task 2 response to an event on IDLE CPU

- CPU state
- Delays

Task 2 gets a chance to run

- CPU is idle
- CPU running

Hardware delays: Wake up from IDLE delay

ISR execution delays

Task response delays

Time

Interrupt 2 Signal
CPU receives signal from IRQ controller
CPU running
Task 2 is woken (placed on CPU run queue)
Response complete

Total Response Delay
Key concepts - interrupt handling (skip)

- Interrupt controller (GIC, APIC etc) sends a hardware signal
- Processor switches mode, banking registers and disabling irq
- Generic Interrupt vector code is called
- Saves the context of the interrupted activity (any context not saved by HW)
- Identify which interrupt occurred, calls relevant ISR

Return:

- Restore context
- Switch processor back to the mode before being interrupted
- Reenable interrupts
Interrupt handling - no non-threaded IRQ nesting

Tglx removes interrupt nesting officially in hard IRQ handlers..

“The following patch series removes the IRQF_DISABLED functionality from the core interrupt code and runs all interrupt handlers with interrupts disabled.”

http://lwn.net/Articles/380536/

Why? Stack overflows is one reason.

Note: In RT kernels, most IRQ handlers so they can preempt each other (these are called secondary handlers with a primary counterpart that wakes them up). What I’m referring in this slide is the primary handler and not any secondary threaded handler.
Kernel preemption is a method used mainly in monolithic and hybrid kernels where all or most device drivers are run in kernel space, whereby the scheduler is permitted to forcibly perform a context switch (i.e. preemptively schedule; on behalf of a runnable and higher priority process) on a driver or other part of the kernel during its execution, rather than co-operatively waiting for the driver or kernel function (such as a system call) to complete its execution and return control of the processor to the scheduler.
Key concepts - kernel preemption (skip)

Cases of Kernel Preemption:
- High priority task wakes up as a result of an interrupt
- Time slice expiration
- System call results in task sleeping

No preemption happens when:
- Preemption explicitly disabled
- Interrupts explicitly disabled
- Spinlock critical sections unless using PREEMPT_RT patchset
- raw spinlock critical sections
RT scheduler team:

"Why did your driver disable preemption?"

Key concepts - long preempt_disable is bad for RT
Key concepts - long preempt_disable is bad for RT

Preemption is disabled after acquiring spinlock (after raw_spin_lock if RT patchset is used)

Preemption off for long time is a problem (high prio task cannot run)

- Use RT patchset for solving lock preempt-off issues, more on that next slide.
- If you have to disable preemption, use need_resched() to check if higher prio task needs CPU to break out of preempt off section.
Key concepts - spinlocks and mutexes (skip)

**Spinlocks - spin till you get the lock**

Good for small critical sections, sections where you can’t sleep.

Preemption is disabled after spinlock is acquired (unless RT patchset is used)

**Mutexes - sleep if you cannot get the lock**

Good for sections where you can sleep.

Critical sections are preemptible

Don’t need to spin and waste CPU (caveat: Linux mutex uses OSQ lock which will spin in some cases, with RT patchset though: mutex lock uses rtmutex)
Key concepts - RT priority inversion (skip)

Issue: a process with more priority can preempt a process holding the lock. The top priority process could wait for a very long time. This happened with standard Linux before version 2.6.18.
Key concepts - RT priority inversion (skip)

Solution: *priority inheritance*. The process holding the lock inherits the priority of the process waiting for the lock with the greatest priority.
Key concepts - RT priority inversion (skip)

rtmutex, spin_lock, mutex code with CONFIG_PREEMPT_RT_FULL have priority-inheritance capabilities.
RT Patchset - rt.wiki.kernel.org

- Spinlock API uses rt-spin_lock - which sleeps while spinning
- Spinlock critical sections are preemptible
- Mutex uses rtmutex which has PI support and doesn’t optimistically spin
- Convert IRQ top halves to use IRQ threads

This talk is not about RT Patchset and its features!
Key concepts - scheduler behavior (skip)

Scheduler needs to put woken up task on CPU, otherwise we’ve latency

Things preventing that:

● Process priority:
  ○ Low prio task waits on the rq while high prio given cpu
● Process scheduling class:
  ○ task is in scheduling class like SCHED_OTHER instead of SCHED_FIFO
● SCHED_FIFO and SCHED_RR always scheduled before SCHED_OTHER/SCHED_BATCH

Note that IRQ threads are SCHED_FIFO tasks with priority 50. Being threads their priority can be changed so that other RT tasks have higher priority.
Real time issues - Categories

● Kernel issues
  ○ Too much time spent in kernel mode
  ○ Preemption turned off
  ○ IRQs turned off
  ○ Spinlocks used where not necessary etc

● Application issues
  ○ App takes too much time in its period
  ○ Compiler issues: suboptimal code
  ○ Poor design - eg. lack of parallel code, cache misses
  ○ CPU frequency
  ○ Wrong scheduling priority, policy
  ○ Page faults in timing critical code

● Hardware issues
  ○ Bus accesses or interconnect take too long
  ○ CPU takes too long to come out of idle
Real time issues - kernel - irqs and preempt off

- Interrupts disabled on local CPU for too long
  - Has the effect of locking the CPU to tasks and ISRs
  - An interrupt that wakes up a task can’t execute its ISR to wake it up

- Preemption disabled on local CPU for too long
  - Has the effect of locking CPU to other tasks
  - Acceptable if preempt off section checks need_resched()
One does not simply disable interrupts!
Real time issues - kernel - irqs off examples

Ex1: (raw_)spinlock_irq_save used where not necessary

Output from irqsoff tracer:

=> started at: atomisp_css2_hw_load
=> ended at: atomisp_css2_hw_load

=> _raw_spin_unlock_irqrestore
=> atomisp_css2_hw_load
=> ia_css_device_load
=> sp_dmem_load
=> ia_css_pipeline_has_stopped
=> ia_css_stream_has_stopped
=> __destroy_stream.isra.4
=> __destroy_streams.constprop.13
=> atomisp_css_stop
Real time issues - kernel - irqs off examples

Ex1: (raw_)spin_lock_irqsave used where not necessary
Real time issues - kernel - irqs off examples

Ex1: (raw_)spin_lock_irqsave used where not necessary

```c
static void atomisp_css2_hw_load(hrt_address addr, void *to, uint32_t n)
{
    unsigned long flags;
    unsigned i;
    char *to = (char *)to;
    unsigned int _from = (unsigned int)addr;

    spin_lock_irqsave(&mmio_lock, flags);
    // replace with:
    // disable_irq_nosync(irq)
    // spin_lock(&mmio_lock)
    raw_spin_lock(&pci_config_lock);
    for (i = 0; i < n; i++, _to++, _from++)
        *to = _hrt_master_port_load_8(_from);
    raw_spin_unlock(&pci_config_lock);
    spin_unlock_irqrestore(&mmio_lock, flags);
}
```
Real time issues - kernel - irqs off examples

Ex1: (raw_)spin_lock_irqsave used where not necessary

Notes:

● With CONFIG_PREEMPT_RT_FULL, spin_lock_irqsave doesn’t disable interrupts. API name maintained for non-RT kernel cases. raw_spin_lock_irqsave still does

● With RT patchset, interrupts run as threads, so irqsave* spinlock variants may not be needed.
Real time issues - kernel - irqs off examples

Ex2: top half (non-threaded) handlers taking a long time

Linux doesn’t support non-threaded IRQ nesting, local interrupts are off when executing a non-threaded interrupt handler. This is also traditionally known as a top half.
Example 2: SST irq top half takes too long. This steals CPU time from the SST thread and other tasks in the system.

Here's function graph tracer output showing top half irq time issue (intel_sst_interrupt_mrfld took 3ms)
Real time issues - kernel - irqs off examples

Ex2: top half handlers taking a long time

Turns out SST was having a hard time accessing the bus… more on that later.

After using function graph tracer to narrow down, I used timestamps + instrumentation to learn that the PCI register space was causing this function to take a long time.

```c
memcpy_fromio(&fw_tstamp, ((void *)(sst_drv_ctx->mailbox + sst_drv_ctx->tstamp) + (str_id * sizeof(fw_tstamp))), sizeof(fw_tstamp));
```

Had to disable PCIe power management features to fix these - more on that later.
Real time issues - kernel - irqs off examples

Ex2: top half handlers taking a long time

Tricks to find top half latency issues

Trick 1: Use function_graph tracer with thresholds set to say 1ms, and depth to 3 on the graph function handle_irq_event

Output on next slide
Real time issues - kernel - irqs off examples

Using ftrace graph to easily find top half handlers taking forever

echo 3 > /d/tracing/max_graph_depth
echo 1000 > /d/tracing/tracing_thresh
trace-cmd record -p function_graph -g handle_irq_event

cat /d/tracing/trace_pipe

0) ! 3165.056 us |      } /* i2c_dw_isr */
0) ! 3168.043 us |    } /* handle_irq_event_percpu */
0) ! 3170.017 us |  } /* handle_irq_event */
0) ! 2757.569 us |      } /* cherryview_irq_handler */
0) ! 2763.906 us |    } /* handle_irq_event_percpu */
0) ! 2766.343 us |  } /* handle_irq_event */
0) ! 3188.289 us |      } /* atomisp_isr [atomisp_css2401a0_v21] */
0) ! 3209.286 us |    } /* handle_irq_event_percpu */
0) ! 3214.460 us |  } /* handle_irq_event */
Real time issues - kernel - irqs off examples

Using ftrace graph to easily find top half handlers taking forever

Trick 2: Use kretprobes

Idea:

1. Install a dynamic probe at handle_irq_event function entry
2. At the entry probe handler, get the IRQ name and timestamp
3. At the exit probe handler, get another timestamp
4. Find entry/exit time difference and print warning if too high
Real time issues - kernel - irqs off examples

Using ftrace graph to easily find top half handlers taking forever

Trick 2: Use kretprobes

/* the entry_handler to timestamp function entry */
static int entry_handler(struct kretprobe_instance *ri, struct pt_regs *regs)
{
    struct irqprobe_data *data;
    struct irq_desc *desc;
    char *str;

    data = (struct irqprobe_data *)ri->data;
    data->entry_stamp = ktime_get();
    str = data->funcname;

    desc = PT_REGS_PARM1(regs);
    str[0] = 0;

    if (desc->action)
        sprint_symbol_no_offset(str, (unsigned long)desc->action->handler);

    return 0;
}
Real time issues - kernel - irqs off examples

Using ftrace graph to easily find top half handlers taking forever

**Trick 2: Use kretprobes**

```c
static int ret_handler(struct kretprobe_instance *ri, struct pt_regs *regs)
{
    int retval = regs_return_value(regs);
    struct irqprobe_data *data = (struct irqprobe_data *)ri->data;
    s64 delta;
    ktime_t now;

    now = ktime_get();
    delta = ktime_to_ns(ktime_sub(now, data->entry_stamp));
    if (delta > 1000 * 1)
        pr_err("IRQ: %s took %lld ns to execute\n",
               data->funcname,
               (long long)delta);
    return 0;
}
```
Real time issues - kernel - irqs off examples

Using ftrace graph to easily find top half handlers taking forever

**Trick 2: Use kretprobes**

```c
static char func_name[NAME_MAX] = "handle_irq_event";
static struct kretprobe irq_kretprobe = {
    .handler = ret_handler,
    .entry_handler = entry_handler,
    .data_size = sizeof(struct irqprobe_data),
    /* Probe up to 20 instances concurrently. */
    .maxactive = 20,
};

static int __init kretprobe_init(void)
{
    int ret;

    irq_kretprobe.kp.symbol_name = func_name;
    ret = register_kretprobe(&irq_kretprobe);
    if (ret < 0) {
        printk(KERN_INFO "register_kretprobe failed, returned %d\n",
               ret);
        return -1;
    }
}
```
Real time issues - kernel - irqs off examples

Using ftrace graph to easily find top half handlers taking forever

Trick 2: Use kretprobes

$ # threshold set to 1ms
$ insmod /lib/modules/thardirq.ko

[ 1002.153168] IRQ: i2c_dw_isr took 3062713 ns to execute
[ 1002.183965] IRQ: i2c_dw_isr took 3158637 ns to execute
[ 1002.202206] IRQ: i2c_dw_isr took 3188238 ns to execute
[ 1002.656567] IRQ: i2c_dw_isr took 3176875 ns to execute
[ 1002.854593] IRQ: i2c_dw_isr took 3161238 ns to execute
[ 1003.157660] IRQ: i2c_dw_isr took 3024987 ns to execute
[ 1003.253201] IRQ: i2c_dw_isr took 3044400 ns to execute
[ 1082.237671] IRQ: sdhci_irq took 1209488 ns to execute
[ 1082.253001] IRQ: sdhci_irq took 1229225 ns to execute
[ 1082.274314] IRQ: sdhci_irq took 1229712 ns to execute
[ 1082.302393] IRQ: i2c_dw_isr took 3167850 ns to execute
[ 1213.410491] IRQ: dhdpcie_isr [bcmdhd] took 533287 ns to execute
Real time issues - kernel - irqs off examples

Ex2: top half handlers taking a long time

Recommendations:

- Use Threaded IRQ
- Time your hard IRQ sections and make sure they’re tiny
Real time issues - kernel - irqs off examples

Ex3: serial console prints in 8250 driver

Serial console prints.. Timestamps show 10ms per print (code execution time between prints was minimal)

```
[  89.966248] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
[  89.975753] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
[  89.985184] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
[  89.994879] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
[  90.004302] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
[  90.013844] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
[  90.023419] atomisp-css2401a0_v21 0000:00:03.0: atomisp_css_isr_thread:no subdev.event:4096
```

Code publicly at ZenFone sources:
Real time issues - kernel - irqs off examples

Ex3: serial console prints in 8250 driver
Serial console prints disable interrupts for a long time (seen upto 6ms per line)

static void
serial8250_console_write(struct console *co, const char *s, unsigned int count)
{
    struct uart_8250_port *up = &serial8250_ports[co->index];
    struct uart_port *port = &up->port;
    unsigned long flags;
    unsigned int ier;
    int locked = 1;

    touch_nmi_watchdog();

    ---> local_irq_save(flags);
    if (port->sysrq) {

Real time issues - kernel - irqs off examples
Ex3: serial console prints in 8250 driver
Serial console prints disable interrupts for a long time

Possible solutions:

● Fix the errors/warning (Usually messages are result of errors/warnings)
● Play with the log levels
  ○ Reduce the log level of the message (use pr_info instead of pr_err)
  ○ Increase printk log level (echo <level> > /proc/sys/kernel/printk)
● Disable serial console - in our final product we disabled this
● Upgrade kernel and use PREEMPT_RT_FULL
Real time issues - kernel - irqs off examples

Ex3: serial console prints in 8250 driver

Note:
Ingo Molnar fixed this already in upstream for -rt

Upstream use spin_lock_irqsave instead of local_irq_save, which ends up not disabling interrupts for CONFIG_PREEMPT_RT_FULL kernels. So if you’re using a fairly recent kernel and have PREEMPT_RT_FULL, you shouldn’t have this problem.
preemptirqsoff tracer

- Start tracing at start of critical section (preempt disabled or irqs off)
- Stop tracing at stop of critical section (preempt enabled and irqs on)
- Show trace with maximum latency
- Can enable function tracing (default on) to show which function executed in critical section

More info: Documentation/trace/ftrace.txt
Real time issues - kernel - real preemptoff issue

Ex4: Lazy max pages

699 static void free_vmap_area_noflush(struct vmap_area *va)
700 {
701     int nr_lazy;
702
703     nr_lazy = atomic_add_return((va->va_end - va->va_start) >> PAGE_SHIFT,
704                                 &vmap_lazy_nr);
705
706     /* After this point, we may free va at any time */
707     llist_add(&va->purge_list, &vmap_purge_list);
708
709     if (unlikely(nr_lazy > lazy_max_pages()))
710         try_purge_vmap_area_lazy();
711 }
Fx4: Lazy max pages

```c
/* Purges all lazily-freed vmap areas.
 *
 * If sync is 0 then don't purge if there is already a purge in progress.
 * If force_flush is 1, then flush kernel TLBs between *start and *end even
 * if we found no lazy vmap areas to unmap (callers can use this to optimise
 * their own TLB flushing).
 * Returns with *start = min(*start, lowest purged address)
 *                  *end = max(*end, highest purged address)
 */
static void __purge_vmap_area_lazy(unsigned long *start, unsigned long *end, int sync, int force_flush)
{
    spin_lock(&purge_lock);

    if (nr) {
        spin_lock(&vmap_area_lock);
        llist_for_each_entry_safe(va, n_va, valist, purge_list)
            __free_vmap_area(va);
        spin_unlock(&vmap_area_lock);
    }
    spin_unlock(&purge_lock);
}
```
Real time issues - kernel - real preemptoff issue

Ex4: Lazy max pages

mm/vmalloc.c (line 593)

593 static unsigned long lazy_max_pages(void)
594 {
595     unsigned int log;
596
597     log = fls(num_online_cpus());
598
599     return log * (32UL * 1024 * 1024 / PAGE_SIZE);
600 }
Real time issues - kernel - real preemptoff issue

Ex4: Lazy max pages. Preemptirqsoff tracer output (with my tracer fix)

# tracer: preemptirqsoff
#
# preemptirqsoff latency trace v1.1.5 on 3.14.37-x86_64-00190-gddfae4b-dirty
# --------------------------------------------------------------------
# latency: 14707 us, #38619/38619, CPU#2 | (M:preempt VP:0, KP:0, SP:0 HP:0 #P:4)
#    -----------------
#    | task: netd-4462 (uid:0 nice:0 policy:0 rt_prio:0)
#    -----------------
#  => started at: __purge_vmap_area_lazy
#  => ended at:   __purge_vmap_area_lazy
#
# (...)-4462 2...1 0us : _raw_spin_trylock <-__purge_vmap_area_lazy
Real time issues - kernel - real preemptoff issue

Ex4: Lazy max pages. Preemptirqsoff tracer output (with my tracer fix)

```c
<...>-4462 2...3 944us : csd_lock_wait.isra.4 <-smp_call_function_many
<...>-4462 2...3 945us+ : csd_lock_wait.isra.4 <-smp_call_function_many
<...>-4462 2...3 948us : preempt_count_sub <-smp_call_function
<...>-4462 2d..2 948us : do_flush_tlb_all <-on_each_cpu
<...>-4462 2...2 949us : preempt_count_sub <-on_each_cpu
<...>-4462 2...1 950us : __raw_spin_lock <-__purge_vmap_area_lazy
<...>-4462 2...1 950us+ : preempt_count_add <-__raw_spin_lock
<...>-4462 2...2 952us+ : __free_vmap_area <-__purge_vmap_area_lazy
<...>-4462 2...2 954us : kfree_call_rcu <-__free_vmap_area
<...>-4462 2...2 955us : __call_rcu.constprop.63 <-kfree_call_rcu
<...>-4462 2d..2 956us : preempt_count_add <-rcu_is_watching
<...>-4462 2d..3 956us : preempt_count_sub <-rcu_is_watching
<...>-4462 2...2 957us : __free_vmap_area <-__purge_vmap_area_lazy

... rinse repeat..
```
RT patchset : s/spin_lock_irqsave/spin_lock/

From include/linux/spin_lock.h

+ifdef CONFIG_PREEMPT_RT_FULL
+  #include <linux/spinlock_rt.h>
+else /* PREEMPT_RT_FULL */
+

From include/linux/spinlock_rt.h

+define spin_lock_irqsave(lock, flags)
+  do {
+    typecheck(unsigned long, flags);
+    flags = 0;
+    spin_lock(lock);
+  } while (0)
RT patchset: s/spin_lock/rt_spin_lock/

From include/linux/spinlock_rt.h

```c
#define spin_lock(lock) do { 
    migrate_disable(); 
    rt_spin_lock(lock); 
} while (0)
```
Real time issues - Hardware: Bus related

- Posted transactions: DONT WAIT for transaction to complete
- Non-Posted transactions: WAIT for transaction to complete.
Real time issues - Hardware: Bus related

Source:
RT Hardware issues

Intel Atom SoC architecture:

- Shared interconnect between peripherals
- Ordering constraints of Accesses on the interconnect
- Behavior of PCIe power sub states
Real time issues - Hardware: : Bus related (wifi recovery)

Broadcom wireless driver sources publicly available at: goo.gl/z1EnJB

Function graph tracer with max depth of 1 ...
# tracer: function_graph
#
# CPU DURATION FUNCTION CALLS
# | | | | | | | |
1) ! 3145.770 us | dhd_bus_cmn_readshared [bcmdhd]()
1) ! 3159.381 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 6.199 us | dhd_bus_cmn_readshared [bcmdhd]()
1) ! 3164.968 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 1.775 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 3.124 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 3.237 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 3.149 us | dhd_bus_cmn_readshared [bcmdhd]()
1) + 40.293 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 3.561 us | dhd_bus_cmn_readshared [bcmdhd]()
1) 3.162 us | dhd_bus_cmn_readshared [bcmdhd]()
Real time issues - Hardware : Bus related (after fixing...)

Broadcom wireless driver sources publicly available at: [goo.gl/z1EnJB](https://goo.gl/z1EnJB)

Function graph tracer with max depth of 1 ...

```plaintext
# tracer: function_graph
#
# CPU DURATION FUNCTION CALLS
# | | | | | | | |
0) + 55.129 us | dhd_bus_cmn_readshared [bcmdhd]()
0)  2.087 us  | dhd_bus_cmn_readshared [bcmdhd]()
0)  3.774 us  | dhd_bus_cmn_readshared [bcmdhd]()
0)  3.949 us  | dhd_bus_cmn_readshared [bcmdhd]()
0) + 55.678 us | dhd_bus_cmn_readshared [bcmdhd]()
0)  1.837 us  | dhd_bus_cmn_readshared [bcmdhd]()
0)  3.612 us  | dhd_bus_cmn_readshared [bcmdhd]()
0)  3.987 us  | dhd_bus_cmn_readshared [bcmdhd]()
0) + 55.504 us | dhd_bus_cmn_readshared [bcmdhd]()
0) + 53.005 us | dhd_bus_cmn_readshared [bcmdhd]()
0)  3.312 us  | dhd_bus_cmn_readshared [bcmdhd]();
```
Others peripheral suffer too, here we see audio was suffering from the same issue:

```c
memcpy_fromio(msg->mailbox_data,
    drv->mailbox + drv->mailbox_recv_offset, size);
```

table:

<table>
<thead>
<tr>
<th>Function Entry</th>
<th>Function Exit</th>
<th>Time, us</th>
</tr>
</thead>
<tbody>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.100</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.975</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.100</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.975</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.362</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.113</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>2.450</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>4.562</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.100</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.900</td>
</tr>
<tr>
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<td>funcgraph_entry</td>
<td>0.087</td>
</tr>
<tr>
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<td>funcgraph_exit</td>
<td>0.925</td>
</tr>
<tr>
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<td>funcgraph_exit</td>
<td>0.100</td>
</tr>
<tr>
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<td>funcgraph_exit</td>
<td>0.100</td>
</tr>
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<td>funcgraph_exit</td>
<td>1.013</td>
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</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_entry</td>
<td>0.900</td>
</tr>
<tr>
<td>funcgraph_entry</td>
<td>funcgraph_exit</td>
<td>9.936</td>
</tr>
<tr>
<td>funcgraph_exit</td>
<td></td>
<td>3167.993</td>
</tr>
</tbody>
</table>

```c
handle_irq_event() {
    _raw_spin_unlock() {
        preempt_count_sub();
    }
}
```
Real time issues - Hardware : CPU wakeup from idle

- PM QoS framework
- \texttt{cpuidle_latency_notify} is called when latency requirement change.
- All cores have to be woken up to calculate new C-state
- Involves sending an IPI (inter-processor interrupt) to all cores to wake-up
- Preemption is turned off until all CPUs wakeup
Real time issues - Hardware : CPU wakeup from idle

/*
 * This function gets called when a part of the kernel has a new latency
 * requirement.  This means we need to get all processors out of their C-state,
 * and then recalculate a new suitable C-state. Just do a cross-cpu IPI; that
 * wakes them all right up.
 */

static int cpuidle_latency_notify(struct notifier_block *b,
        unsigned long l, void *v)
{
    smp_call_function(smp_callback, NULL, 1);
    return NOTIFY_OK;
}

int smp_call_function(smp_call_func_t func, void *info, int wait)
{
    preempt_disable();
    smp_call_function_many(cpu_online_mask, func, info, t);
    preempt_enable();

    return 0;
}
Real time issues - Hardware : CPU wakeup from idle

mmcqd/0-1408 0...1 0us : smp_call_function <-cpuidle_latency_notify
mmcqd/0-1408 0...1 1us : smp_call_function_many <-smp_call_function
mmcqd/0-1408 0...1 2us : csd_lock_wait.isra.4 <-smp_call_function_many
mmcqd/0-1408 0...1 3us : csd_lock_wait.isra.4 <-smp_call_function_many
mmcqd/0-1408 0...1 3us : csd_lock_wait.isra.4 <-smp_call_function_many
mmcqd/0-1408 0...1 4us : native_send_call_func_ipi <-smp_call_function_many
mmcqd/0-1408 0...1 5us : flat_send_IPI_allbutself <-native_send_call_func_ipi ← Ensure its unlocked
mmcqd/0-1408 0...1 6us+ : csd_lock_wait.isra.4 <-smp_call_function_many ← Send IPI to all
mmcqd/0-1408 0...1 10us! : csd_lock_wait.isra.4 <-smp_call_function_many ← first csd_lock_wait succeeds
mmcqd/0-1408 0d..1 121us : smp_apic_timer_interrupt <-apic_timer_interrupt ← second one always delayed
mmcqd/0-1408 0d..1 122us : irq_enter <-smp_apic_timer_interrupt

...
Real time issues - Hardware : CPU wakeup from idle

Late wake up on Core 3 causing preempt off delay on Core 0
Real time - application issues

Fig. (1a) - Example of working periodic real-time tasks with deadlines equal to periods. Not deadline misses.

Fig. (1b) - Example of problematic periodic real-time task with deadline misses due to too long execution time.

Fig. (1c) - Example of problematic periodic real-time task with deadline misses due to too jitter (scheduling latency).

Fig. (1d) - Example of problematic periodic real-time task with deadline misses due to too jitter combined with too long execution time.

Image Source: Eduardo Valentin <edubesval@gmail.com>
Amazon Echo

- A product that has stringent requirements on response time and customer experience
- Always listening for “Alexa”
- Examples of audio algorithms always running.
  - From amazon.com/echo product page
    - “Tucked under the light ring is an array of seven microphones that use beam-forming technology and enhanced noise cancellation. With far-field voice recognition, Echo can hear you ask a question from any direction—even while playing music.”
Analysis of Audio pipeline with Android systrace
Real time issues: application - CPU frequency

<table>
<thead>
<tr>
<th>Kernel</th>
<th>1200 MHz</th>
<th>600 MHz</th>
<th>1800 MHz</th>
</tr>
</thead>
<tbody>
<tr>
<td>CPU 0: Clock Frequency:</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CPU 1: Clock Frequency:</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CPU 2: Clock Frequency:</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>CPU 3: Clock Frequency:</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Process 2433</th>
<th>AudioIn_15</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>A</td>
</tr>
<tr>
<td></td>
<td>AL</td>
</tr>
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<td></td>
<td>ALG</td>
</tr>
<tr>
<td></td>
<td>IN</td>
</tr>
<tr>
<td></td>
<td>ALG</td>
</tr>
<tr>
<td></td>
<td>AL</td>
</tr>
<tr>
<td></td>
<td>A</td>
</tr>
</tbody>
</table>

Things start falling apart
Real time issues: application - Cache misses

Analysis of Audio bottlenecks with perf utility

```
# To display the perf.data header info, please use --header/--header-only options.
#
# Samples: 4K of event 'cache-misses'
# Event count (approx.): 11166978
#
# Overhead Command Shared Object
# .......................... ...........................
# 30.21% mediaserver libc.so [] memmove
 # memmove
--80.10%--

Audio Algorithm Hotpath

_ZN7android6Thread11_threadLoopEPv
_ZN11thread_data_t10_trampolineEPKs
_ZL15_pthread_startPv
_start_thread
__bionic_clone
0x0
_ZN7android6Parcel14releaseObjectsEv
_ZN7android6Parcel14freeDataNoInitEv
_ZN7android6ParcelD1Ev
_ZN7android20BpAudioFlingerClient15ioConfigChangedEiiPKv
0x0
0x648d0006
```

Generic Android Audio Pipeline
Real time issues: application - Other issues?

- Lack of parallelization
- Compiler issues
- Memory locality
- Page Faults
Real time issues: system - Scheduling

- Get your priorities right
- Use the right policy
Real time issues: Scheduling: Find issues..

Using scheduler delay statistics (CONFIG_SCHEDSTATS)

cat /proc/<pid>/sched for AudioIn thread in Android (CFS policy)

```
se.exec_start : 115703.404109
se.vruntime : 1761777.444578
se.sum_exec_runtime : 1619.691347
se.statistics.wait_start : 0.000000
se.statistics.sleep_start : 118028.360499
se.statistics.block_start : 0.000000
se.statistics.sleep_max : 23528.490473
se.statistics.block_max : 0.009298
se.statistics.exec_max : 6.342921 <- only for high load tasks
se.statistics.slice_max : 0.430483
se.statistics.wait_max : 12.998756
se.statistics.wait_sum : 279.959758
```
Real time issues: Scheduling: Find issues..

cat /proc/<pid>/sched for AudioIn thread in Android (RT policy)

pid 4732's current scheduling policy: SCHED_FIFO
pid 4732's current scheduling priority: 2

AudioIn_1A (4732, #threads: 34)

<table>
<thead>
<tr>
<th>Variable</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>se.exec_start</td>
<td>287659.027531</td>
</tr>
<tr>
<td>se.vruntime</td>
<td>-4.755003</td>
</tr>
<tr>
<td>se.sum_exec_runtime</td>
<td>11894.018812</td>
</tr>
<tr>
<td>se.statistics.wait_start</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.sleep_start</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.block_start</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.sleep_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.block_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.exec_max</td>
<td>6.541440</td>
</tr>
<tr>
<td>se.statistics.slice_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.wait_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.wait_sum</td>
<td>0.000000</td>
</tr>
</tbody>
</table>

Not that useful for RT!
Real time issues: Scheduling: Find issues..

Getting scheduling delays for any scheduler policy

/proc/pid/schedstat already calculates total run queue delays per task

Why not also find the maximum run delay?

```c
static inline void sched_info_dequeued(struct rq *rq, struct task_struct *t)
    delta = now - t->sched_info.last_queued;
    sched_info_reset_dequeued(t);
    t->sched_info.run_delay += delta;
+       schedstat_set(t->se.statistics.run_delay_max,
+                      max(delta, t->se.statistics.run_delay_max));
```
Real time issues: Scheduling: Find issues..

maximum runqueue delay in `/proc/<pid>/schedstat` for RT task

<table>
<thead>
<tr>
<th></th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>se.exec_start</td>
<td>495117.327836</td>
</tr>
<tr>
<td>se.vruntime</td>
<td>-5.000000</td>
</tr>
<tr>
<td>se.sum_exec_runtime</td>
<td>22153.249640</td>
</tr>
<tr>
<td>se.statistics.wait_start</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.sleep_start</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.block_start</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.sleep_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.block_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.exec_max</td>
<td>0.881424</td>
</tr>
<tr>
<td>se.statistics.exec_hist[HIST_0_10US]</td>
<td>5864</td>
</tr>
<tr>
<td>se.statistics.exec_hist[HIST_10US_100US]</td>
<td>4834</td>
</tr>
<tr>
<td>se.statistics.exec_hist[HIST_100US_1MS]</td>
<td>22478</td>
</tr>
<tr>
<td>se.statistics.exec_hist[HIST_1MS_10MS]</td>
<td>0</td>
</tr>
<tr>
<td>se.statistics.exec_hist[HIST_10MS_100MS]</td>
<td>0</td>
</tr>
<tr>
<td>se.statistics.exec_hist[HIST_MORE_100MS]</td>
<td>0</td>
</tr>
<tr>
<td>se.statistics.slice_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.wait_max</td>
<td>0.000000</td>
</tr>
<tr>
<td>se.statistics.run_delay_max</td>
<td>0.377203</td>
</tr>
</tbody>
</table>
Real time issues: Scheduling: Find issues..

Room for improvement even with this:

- We handle cases where a task is queued -> CPU

- What about cases where a task is migrated after queuing? (queued -> dequeued -> queued on another rq -> CPU)

- This works for run_delay because its cumulative, but not so much for run_delay_max:

Here’s a modification of Steven’s rt lock test (https://lwn.netArticles/425583/)

<table>
<thead>
<tr>
<th>Task</th>
<th>run_delay_max</th>
<th>rdm_naive</th>
<th>run_migr_max</th>
<th>nr_running_migr</th>
<th>old_rq_delay_max</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>8.955969</td>
<td>8.955969</td>
<td>0.000947</td>
<td>449</td>
<td>5.039939</td>
</tr>
<tr>
<td>1</td>
<td>8.494575</td>
<td>5.974765</td>
<td>0.000779</td>
<td>677</td>
<td>4.049992</td>
</tr>
<tr>
<td>2</td>
<td>8.462174</td>
<td>5.958943</td>
<td>0.010491</td>
<td>572</td>
<td>3.048980</td>
</tr>
<tr>
<td>3</td>
<td>5.514705</td>
<td>5.514705</td>
<td>0.000396</td>
<td>3</td>
<td>0.000579</td>
</tr>
<tr>
<td>4</td>
<td>5.557162</td>
<td>5.548496</td>
<td>0.000747</td>
<td>6</td>
<td>5.548496</td>
</tr>
<tr>
<td>5</td>
<td>5.953767</td>
<td>5.953767</td>
<td>0.000410</td>
<td>2</td>
<td>0.008089</td>
</tr>
<tr>
<td>6</td>
<td>0.031044</td>
<td>0.027696</td>
<td>0.000281</td>
<td>1</td>
<td>0.003067</td>
</tr>
<tr>
<td>7</td>
<td>0.003993</td>
<td>0.003993</td>
<td>0.000000</td>
<td>0</td>
<td>0.000000</td>
</tr>
</tbody>
</table>
Real time issues: Scheduling: Find issues..

Cyclic test

“cyclic test measures the delta from when it's scheduled to wake up from when it actually does wake up”

Clark Williams, An Overview of Realtime Linux, Redhat Summit, June 18-20th, 2008.

https://rt.wiki.kernel.org/index.php/Cyclic_test
Real time issues: Tools: latency_hist (demo)

Latency Hists (available in RT Patchset)
CONFIG_INTERRUPT_OFF_LATENCY
CONFIG_PREEMPT_OFF_LATENCY
CONFIG_WAKEUP_LATENCY

• Possible latencies:
  ○ Preemption Off histogram
  ○ IRQs Off histogram

• Effective latencies:
  ○ Histogram of wake up latency per CPU
  ○ Details of Task experiencing the latency
Real time issues: Tools: latency_hist

- Demo:
  - Example RT task: Cyclic test
  - Example kernel module: introduce Preempt Off latency
- Results:
  - Cyclic test shows latency
  - Latency hists shows hists & maximum effective latency
Real time issues: Tools: latency_hist

Code in ‘trouble maker’ kernel module:

```c
int x;

static int __init test_module_init(void)
{
    unsigned long j, i, loop1 = 100, loop2 = 10000;

    /* Introduces a preempt delay about about 50ms */
    preempt_disable();

    for (i = 0; i < loop1 * loop2; i++)
        ACCESS_ONCE(x) += 1;

    preempt_enable();

    return -1;
}
```
Real time issues: Tools: latency_hist

Running trouble maker module in a loop:

while [ 1 ]; do insmod ./preemptd.ko; done

Run cyclicstest with priority 80:

./cyclicstest -t1 -p 80 -n -i 10000 -l 10000
Real time issues: Tools: latency_hist

cyclicictest gets victimized sooner or later:

root@raspberrypi:/home/pi# ./cyclicictest -t1 -p 80 -n -i 10000 -l 10000
# /dev/cpu_dma_latency set to 0us
policy: fifo: loadavg: 0.18 0.12 0.06 1/125 1382

T: 0 (1075) P:80 I:10000 C: 5252 Min: 14 Act: 36 Avg: 36 Max: 2024
Real time issues: Tools: latency_hist

latency_hist histograms can show per-CPU latency histogram:

```bash
# cat /sys/kernel/debug/latency_hist/wakeup/CPU3
...

583       1
718       1
772       1
807       1
850       1
853       1
861       1
1120      1
1470      1
1895      1
1994      1
```
Real time issues: Tools: latency_hist

latency_hist shows **details of max wakeup latency** per-CPU:

```plaintext
# cat /sys/kernel/debug/latency_hist/wakeup/max_latency-CPU3

CPU 3 max latency info:
1075 80 1994 (0) cyclictest <- 1329 -21 insmod 1107.418695
```
Other tools

Latency tracker
Rt app

Thank you